

MSO definable string transductions and two-way finite state transducers

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Abstract

String transductions that are definable in monadic second-order (mso) logic (without the use of parameters) are exactly those realized by deterministic two-way finite state transducers. Nondeterministic mso definable string transductions (i.e., those definable with the use of parameters) correspond to compositions of two nondeterministic two-way finite state transducers that have the finite visit property. Both families of mso definable string transductions are characterized in terms of Hennie machines, i.e., two-way finite state transducers with the finite visit property that are allowed to rewrite their input tape.

Introduction

In language theory, it is always a pleasant surprise when two formalisms, introduced with different motivations, turn out to be equally powerful, as this indicates that the underlying concept is a natural one. Additionally, this means that notions and tools from one formalism can be made use of within the other, leading to a better understanding of the formalisms under consideration. Most famous in this respect are of course the regular languages [Yu97], that can be defined using a computational formalism (finite state automata, either deterministic or nondeterministic), but also have well-known grammatical (right-linear grammars), operational (rational operations), algebraic (congruences of finite index), and logical (monadic second-order logic

of one successor) characterizations [MCPi43, RaSc59, Cho56, Kle56, Myh57, Ner58, Büc60, Elg61].

In this paper we study ‘regular’ (string-to-string) transductions, rather than regular languages, and we obtain the equivalence of particular computational and logical formalisms, modestly following in the footsteps of Büchi and Elgot. Their original work [Büc60, Elg61], demonstrating how a logical formula may effectively be transformed into a finite state automaton accepting the language specified by the formula when interpreted over finite sequences, shows how to relate the *specification* of a system behaviour (as given by the formula) to a possible *implementation* (as the finite state behaviour of an automaton). In recent years much effort has been put into transforming these initial theoretical results into software tools for the verification of finite state systems, *model checking*, see the monograph [Kur94]. Generalizations of the result of Büchi and Elgot include infinite strings [Büc62], trees [Don70, ThWr68], traces (a syntactic model for concurrency) [Ebi95], texts (strings with an additional ordering) [HoPa97], and tree-to-tree transductions [BlEn97, EnMa98]. We refer to [Tho97] for an overview of the study of formal languages within the framework of mathematical logic.

We give a short description of the two formalisms of ‘regular’ string transductions that we study in this paper. We mainly consider the deterministic case.

A two-way finite state transducer (or two-way generalized sequential machine, 2gsm) is a finite state automaton equipped with a two-way input tape, and a one-way output tape. Such a transducer may freely move over its input tape, and may typically reverse or copy parts of its input string. It is, e.g., straightforward to construct a transducer realizing the relation $\{(w, ww) \mid w \in \{a, b\}^*\}$. It should be clear from this example that regular languages are not closed under 2gsm mappings, contrary to their closure under one-way gsm mappings.

However, it is well known [RaSc59, She59, HoUl79] that two-way finite state automata accept only regular languages, and consequently (using a straightforward direct product construction) the regular languages are closed under inverse 2gsm transductions. From this general result we may infer a large number of specific closure properties of the regular languages, such as closure under the ‘root’ operation $\sqrt{K} = \{w \mid ww \in K\}$. It is maybe less well known that the (deterministic) 2gsm mappings are closed under composition [ChJá77]. This result is used as a powerful tool in this paper.

The monadic second-order (mso) logic of one successor is a logical framework that allows one to specify string properties using quantification over sets of positions in the string. As stated above, Büchi and Elgot proved that the string languages specified by mso definable properties are exactly the regular languages. The logic has a natural generalization to graphs, with quantification over sets of nodes, and predicates referring to node labels and edge labels. It is used to define graph-to-graph transductions, by specifying the edges of the output graph in terms of properties of (copies of) a given input graph [Cou97, Eng97]. This is just a special case of the notion of interpretation of logical structures, well known in mathematical logic (see, e.g., [See92, Section 6]). These mso definable graph transductions play an important role in the theory of graph rewriting, as the two main families of context-free graph languages can be obtained by applying mso definable graph transductions to regular tree languages [EnOo97, CoEn95].

Here we consider mso definable string transductions, i.e., the restriction of mso definable graph transductions to linear input and output graphs. It is known that mso definable (string) transductions are closed under composition, and that the regular languages are closed under inverse mso definable transductions (recall that regular is equivalent to mso definable), see, e.g., [Cou94].

Apart from these similar closure properties there is more evidence in the literature that indicates the close connection between 2gsm transductions and mso definable transductions. First, various specific 2gsm transductions were shown to be mso definable, such as one-way gsm mappings, mirror image, and mapping the string w onto w^n (for fixed n), cf. [Cou97, Prop 5.5.3]. Second, returning to the theory of graph grammars, it is explained in [Eng97, pages 192–8] that the ranges (i.e., output languages) of mso definable (string) transductions are equal to the (string) languages defined by linear context-free graph grammars, which, by a result of [EnHe91], equal the ranges of 2gsm transductions. Consequently, the two families of transductions we consider have the same generative power (on regular input). This, however, does not answer the question whether they are the same family of transductions (cf. Section 6 of [Cou94]). In this paper we answer this question positively (in the deterministic case). Thus, string transductions that are *specified* in mso logic can be *implemented* on 2gsm's, and vice versa.

Our paper is organized as follows.

In a preliminary section we mainly recall notions and notations regarding

graphs, in particular mso logic for graphs and strings. Moreover, we recall the usual, natural representation of strings as linear graphs that allows a transparent interpretation of strings and string languages within the setting of the mso logic for graphs.

In Section 2 we study *two-way machines*, our incarnation of two-way generalized sequential machines. We extend the basic model by allowing the machines to ‘jump’ to new positions on the tape (not necessarily adjacent to the present position) as specified by an mso formula that is part of the instructions. This ‘hybrid’ model (in between logic and machine) facilitates the proof of our main result. We consider yet another variant of the 2gsm which allows ‘regular look-around’, i.e., the ability to test the strings to the left and to the right of the reading head for membership in a regular language. The equivalence of the basic 2gsm model and our two extended models (in the deterministic case) is demonstrated using the closure of 2gsm under composition and using Büchi and Elgot’s result for regular languages.

In Section 3 we recall the definition of mso definable graph transduction, and restrict that general notion to mso definable string transductions by considering graph representations for strings. In addition to the representation of Section 1, we use an alternative, natural and well-known, graph representation for strings. Again it uses linear graphs, with labels on the edges rather than on the nodes to represent the symbols of the string. These two representations differ slightly, due to an unfortunate minor technicality involving the empty string; the second representation gives more uniform results.

The main result of the paper is presented as Theorem 23: the equivalence of the (deterministic) 2gsm from Section 2, and the mso definable string transductions from Section 3. Section 4 contains the proof of this result. In order to transform a 2gsm into the mso formalism we consider the ‘computation space’ of a 2gsm on a given input. This is the graph which has a node for each pair consisting of a tape position and a state of the 2gsm. These nodes are connected by edges representing the possible moves of the 2gsm. The transduction is then decomposed into (basically) two constructions, each of which is shown to be mso definable. First the computation space is defined in terms of the input string, then the computation path for the input (and its resulting output string) is recovered from the computation graph. One implication of the main result then follows by the closure of mso definable (graph!) transductions under composition. The reverse implication is obtained by transforming an mso definable string transduction into a 2gsm equipped with mso instructions, the tool we introduced in Section 2.

In Section 5 we study nondeterminism. This feature can be added to mso definable transductions by introducing so-called ‘parameters’: free set variables in the definition of the transduction [Cou97]. The output of the transduction for a given input may then vary for different valuations of these parameters. These transductions are closed under composition, as opposed to those realized by nondeterministic 2gsm. We conclude that as opposed to the deterministic case, the two nondeterministic families are incomparable. Finally, we observe that the family of nondeterministic mso transductions is equal to the family of transductions defined by composing a (nondeterministic) relabelling and a deterministic transduction.

Finite visit machines form the topic of our final section, Section 6. These machines have a fixed bound on the number of times each of the positions of their input tape may be visited during a computation. We characterize the nondeterministic mso definable string transductions as compositions of two nondeterministic 2gsm’s with the finite visit property. Additionally we demonstrate that an arbitrary composition of nondeterministic 2gsm’s realizes a nondeterministic mso definable string transduction if and only if that transduction is finitary, i.e., it has a finite number of images for every input string.

A more direct characterization can be obtained by considering Hennie transducers, i.e., finite visit 2gsm’s that are allowed to rewrite the symbols on their input tape. These machines characterize the mso definable transductions, both in the deterministic case [ChJá77] and the nondeterministic case.

An extended abstract of this paper is published as [EnHo99].

1 Preliminaries

We recall some notions and results regarding graphs and their monadic second order logic.

By $|w|$ we denote the length of the string w .

We use \circ to denote the composition of binary relations (note the order): $R_1 \circ R_2 = \{(w_1, w_3) \mid \text{there exists } w_2 \text{ such that } (w_1, w_2) \in R_1, (w_2, w_3) \in R_2\}$, and extend it to families of binary relations: $F_1 \circ F_2 = \{R_1 \circ R_2 \mid R_1 \in F_1, R_2 \in F_2\}$.

A binary relation R is *functional*, if $(w, z_1) \in R$ and $(w, z_2) \in R$ imply $z_1 = z_2$. It is *finitary*, if each original is mapped to only finitely many images, i.e., the set $\{z \mid (w, z) \in R\}$ is finite for each w in the domain of R .

Graphs. Let Σ and Γ be alphabets of node labels and edge labels, respectively. A *graph* over Σ and Γ is a triple $g = (V, E, \ell)$, where V is the finite set of nodes, $E \subseteq V \times \Gamma \times V$ the set of edges, and $\ell : V \rightarrow \Sigma$ the node labelling. The set of all graphs over Σ and Γ is denoted by $\text{GR}(\Sigma, \Gamma)$. We allow graphs that have both labelled and unlabelled nodes and edges by introducing a designated symbol $*$ to represent an ‘unlabel’ in our specifications, but we omit this symbol from our drawings. We write $\text{GR}(*, \Gamma)$ and $\text{GR}(\Sigma, *)$ to distinguish the cases when all nodes are unlabelled, and all edges are unlabelled, respectively.

Logic for graphs. For alphabets Σ and Γ , the monadic second-order logic $\text{MSO}(\Sigma, \Gamma)$ expresses properties of graphs over Σ and Γ . The logical language uses both node variables x, y, \dots and node-set variables X, Y, \dots .

There are four types of atomic formulas: $\text{lab}_\sigma(x)$, meaning node x has label σ (with $\sigma \in \Sigma$); $\text{edge}_\gamma(x, y)$, meaning there is an edge from x to y with label γ (with $\gamma \in \Gamma$); $x = y$, meaning nodes x and y are equal; and $x \in X$, meaning x is an element of X .

As usual, formulas are built from atomic formulas with the propositional connectives $\neg, \wedge, \vee, \rightarrow$, using the quantifiers \forall and \exists both for node variables and node-set variables.

A useful example [ThWr68] of such a formula is the binary predicate \preceq claiming the existence of a (directed) path from x to y :

$$x \preceq y = (\forall X)[(x \in X \wedge \text{closed}(X)) \rightarrow y \in X]$$

where $\text{closed}(X) = (\forall z_1)(\forall z_2)(z_1 \in X \wedge \text{edge}(z_1, z_2) \rightarrow z_2 \in X)$, and $\text{edge}(z_1, z_2) = \bigvee_{\gamma \in \Gamma} \text{edge}_\gamma(z_1, z_2)$. We also use $x \prec y$, where one additionally requires that $x \neq y$; for acyclic graphs this expresses the existence of a nonempty path from x to y .

Let φ be a formula of $\text{MSO}(\Sigma, \Gamma)$ with set Ξ of free variables (of either type), and let $g = (V, E, \ell)$ be a graph in $\text{GR}(\Sigma, \Gamma)$. Let ν be a valuation of φ , i.e., a mapping that assigns to each node variable $x \in \Xi$ an element $\nu(x)$ of V , and to each set variable $X \in \Xi$ a subset $\nu(X)$ of V . We write $g, \nu \models \varphi$ if φ is satisfied in the graph g , where the free variables of φ are valued according to ν .

Let $\varphi(x_1, \dots, x_m, X_1, \dots, X_n)$ be an $\text{MSO}(\Sigma, \Gamma)$ formula with free node variables x_i and free node set variables X_j , and let u_1, \dots, u_m be nodes of graph g , and U_1, \dots, U_n sets of nodes of g . We write $g \models \varphi(u_1, \dots, u_m, U_1, \dots, U_n)$ whenever $g, \nu \models \varphi(x_1, \dots, x_m, X_1, \dots, X_n)$, where ν is the valuation with $\nu(x_i) = u_i$, $\nu(X_j) = U_j$.

Let Ξ be a finite set of variables. The set $\{0, 1\}^\Xi$ of 0, 1-assignments to elements of Ξ is finite, and may be considered as an alphabet. A Ξ -valuated graph over Σ and Γ is a graph in $\text{GR}(\Sigma \times \{0, 1\}^\Xi, \Gamma)$, such that for every node variable x in Ξ there is a unique node of the graph of which the label $(\sigma, f) \in \Sigma \times \{0, 1\}^\Xi$ satisfies $f(x) = 1$.

Clearly, such a Ξ -valuated graph g determines a graph $g|_\Sigma$ in $\text{GR}(\Sigma, \Gamma)$, by dropping the $\{0, 1\}^\Xi$ component of its node labels, as well as a valuation ν_g of the variables in Ξ , by taking

- for a node variable $x \in \Xi$, $\nu_g(x) = u$, where u is the unique node having a label (σ, f) with $f(x) = 1$,
- for a node-set variable $X \in \Xi$, $\nu_g(X) = U$, where U consists of all nodes v having a label (σ, f) with $f(X) = 1$.

For a formula φ of $\text{MSO}(\Sigma, \Gamma)$ with free variables in Ξ , and a Ξ -valuated graph g we write $g \models \varphi$ if φ is true for the underlying graph under the implicitly defined valuation, i.e., if $g|_\Sigma, \nu_g \models \varphi$; φ defines the graph language $GL(\varphi) = \{g \in \text{GR}(\Sigma \times \{0, 1\}^\Xi, \Gamma) \mid g \models \varphi\}$. A graph language is *mso definable* if there exists a closed mso formula that defines the language.

String representation. A string $w \in \Sigma^*$ of length k can be represented by the graph $\text{nd-gr}(w)$ in $\text{GR}(\Sigma, *)$, consisting of k nodes labelled by the

consecutive symbols of w , with $k - 1$ (unlabelled) edges representing the successor relation for the positions of the string. In the figure below, we show $\text{nd-gr}(ababb)$. Note that for the empty string λ , $\text{nd-gr}(\lambda)$ is the empty graph. With this representation, a formula φ of $\text{MSO}(\Sigma, *)$ defines the string language $L(\varphi) = \{w \in (\Sigma \times \{0, 1\}^\Xi)^* \mid \text{nd-gr}(w) \models \varphi\}$, where Ξ is the set of free variables of φ ; note that $\text{nd-gr}(w)$ is a Ξ -valuated graph over Σ and $*$.



Given the close connection between the positions and their successor relation in a string w on the one hand, and the nodes and their connecting edges in $\text{nd-gr}(w)$ on the other, we say that a string w satisfies a formula φ if $\text{nd-gr}(w) \models \varphi$.

String languages definable by monadic second-order formulas are exactly the regular languages, as shown by Büchi and Elgot.

1 Proposition ([Büc60, Elg61])

1. $L(\varphi)$ is a regular string language for every formula φ of $\text{MSO}(\Sigma, *)$.
2. A string language $K \subseteq \Sigma^*$ is regular iff there is a closed formula φ of $\text{MSO}(\Sigma, *)$ such that $K = L(\varphi)$.

We will also refer to Proposition 1 as ‘Büchi’s result’, with due apologies to Elgot.

Observe that the set of all strings over a fixed alphabet Σ forms an mso definable *graph* language via the above representation. The defining formula for the set $\{\text{nd-gr}(w) \mid w \in \Sigma^*\}$ over $\text{MSO}(\Sigma, *)$ expresses the existence of an initial and a final node (provided the graph is nonempty) and demands that every node has at most one direct successor (i.e., the edge relation is functional); ‘guards’ $(\exists x)\text{true} \rightarrow$ are added in order to make the empty string λ satisfy the formula.

$$\begin{aligned}
& (\exists x)\text{true} \rightarrow (\exists x)(\forall y)(x \preceq y \wedge \neg(y \prec x)) \\
& \wedge (\exists x)\text{true} \rightarrow (\exists x)(\forall y)(y \preceq x \wedge \neg(x \prec y)) \\
& \wedge (\forall x)(\forall y_1)(\forall y_2)((\text{edge}(x, y_1) \wedge \text{edge}(x, y_2)) \rightarrow y_1 = y_2)
\end{aligned}$$

As a consequence, the set of graphs representing a string language K , $\{\text{nd-gr}(w) \mid w \in K\}$ is an mso definable graph language for every regular language K .

2 Two-Way Machines

We present our (slightly nonstandard) model of two-way generalized sequential machines (2gsm), or two-way finite state transducers. In order to facilitate the proof of the equivalence of two-way finite state transductions and logically definable transductions we extend the basic model to a machine model that has its input tests as well as moves specified by mso formulas. We prove the equivalence of this extended model to the basic model. An important tool in this proof is the observation that a two-way automaton is able to keep track of the state of another (one-way) finite state automaton (proved in Lemma 3 of [HoUl67], see also p. 212 of [AHU69]). We formalize this fact by extending the 2gsm with the feature of ‘regular look-around’. The equivalence of this model with the basic model is then proved using the related result of [ChJá77] stating that deterministic two-way finite state transductions are closed under composition. The equivalence of the regular look-around model with the mso formula model is proved using Büchi’s result (Proposition 1).

Since we need several types of two-way machines, we first introduce a generic model, and then instantiate it in several ways.

A *two-way machine* (2m) is a finite state device equipped with a two-way input tape (read only), and a one-way output tape. In each step of a computation the machine reads an input symbol, changes its internal state, outputs a string, and moves its input head, all depending on the symbol read and the original internal state.

We specify a 2m as a construct $\mathcal{M} = (Q, \Sigma_1, \Sigma_2, \delta, q_{in}, q_f)$, where Q is the finite set of states, Σ_1 and Σ_2 are the input alphabet and output alphabet, q_{in} and q_f are the initial and the final state, and δ is a finite set of *instructions*. Each instruction is of the form $(p, t, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0)$, where $p \in Q - \{q_f\}$ is the present state of the machine, t is a test to be performed on the input, and the triples (q_i, α_i, μ_i) , $i = 1, 0$, fix the action of the machine depending on the outcome of the test t : $q_i \in Q$ is the new state, $\alpha_i \in \Sigma_2^*$ is the string written on the output tape, and μ_i describes the (deterministic) move of the reading head on the input tape. The precise form of these instructions varies from one model to another, in particular the form of the test t , and the moves μ_i .

The above instruction can be expressed as the following informal code.

```

label  $p$ :  if  $t$  then write  $\alpha_1$  ; move  $\mu_1$  ; goto  $q_1$ 
           else write  $\alpha_0$  ; move  $\mu_0$  ; goto  $q_0$ 
fi

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The string on the input tape is marked by two special symbols, \vdash and \dashv , indicating the boundaries of the tape. So, when processing the string $\sigma_1 \cdots \sigma_n$, $\sigma_i \in \Sigma_1$, the tape has $n + 2$ reachable positions $0, 1, \dots, n, n + 1$, containing the string $\vdash \sigma_1 \cdots \sigma_n \dashv$. The reading head is on one of these positions.

The 2m \mathcal{M} realizes the transduction $m \subseteq \Sigma_1^* \times \Sigma_2^*$, such that $(w, z) \in m$ whenever there exists a computation with $\vdash w \dashv$ on the input tape, starting in initial state q_{in} with the input head on position 0 (where the symbol \vdash is stored), and ending in the accepting state q_f , while z has been written on the output tape.

A 2m is *deterministic* if for each state p there is at most one instruction $(p, t, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0)$ that starts in p . Note that the transduction m realized by a deterministic 2m \mathcal{M} is a partial function $m : \Sigma_1^* \rightarrow \Sigma_2^*$ because the μ_i in the instructions describe deterministic moves of the reading head.

We consider the usual two-way *generalized sequential machine* (2gsm), introduced in [AhU170], and two new instantiations of the generic 2m model, the 2gsm with *regular look-around*, and the 2gsm with *mso-instructions*.

2gsm. For the basic 2gsm model each instruction $(p, t, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0)$ in δ satisfies $t \in \Sigma_1 \cup \{\vdash, \dashv\}$, and $\mu_i \in \{-1, 0, +1\}$, $i = 1, 0$.

Executing an instruction $(p, \sigma, q_1, \alpha_1, \epsilon_1, q_0, \alpha_0, \epsilon_0) \in \delta$ the 2gsm, assuming it is in internal state p , when reading σ on its input tape, changes its state to q_1 , writes α_1 to its output tape, and moves its head from the present position i to the position $i + \epsilon_1$ (provided $0 \leq i + \epsilon_1 \leq n + 1$); if σ is not read on the input tape it acts similarly according to the triple $(q_0, \alpha_0, \epsilon_0)$. Recall that there are no instructions starting in the final state.

It is more customary to formalize the instructions of a 2gsm as 5-tuples $(p, \sigma, q, \alpha, \epsilon)$, not having the ‘else-part’ of our instructions. These two approaches are easily seen to be equivalent. Obviously, the 5-tuple can be extended to an 8-tuple by adding a dummy ‘else-part’, as in $(p, \sigma, q, \alpha, \epsilon, p, \lambda, 0)$. Conversely, one of our instructions $(p, \sigma, q_1, \alpha_1, \epsilon_1, q_0, \alpha_0, \epsilon_0)$ can be replaced by the ‘if-part’ $(p, \sigma, q_1, \alpha_1, \epsilon_1)$ and all alternatives $(p, \sigma', q_0, \alpha_0, \epsilon_0)$, $\sigma' \neq \sigma$.

For determinism we require each state to have at most one instruction, whereas the customary notion considers both state and input symbol. This,

somewhat unusual, formulation allows us to have the above common definition of determinism for all necessary instantiations of our generic model, without having to worry about the mutual exclusiveness of the tests t . This is the reason for choosing our 8-tuple formalism.

The first of the two translations (from 5-tuple model to our 8-tuple model) does not respect determinism. We can solve this by checking all alternatives in a given state consecutively, as follows. Let $(p, \sigma_i, q_i, \alpha_i, \epsilon_i)$, $i = 1, \dots, k$ be all the instructions for state p in a deterministic (5-tuple) 2gsm, which means that the σ_i are different. Introduce $k + 1$ copies $p = p^{(1)}, p^{(2)}, \dots, p^{(k)}, p^{(k+1)}$ of p . Then, the instructions $(p^{(i)}, \sigma_i, q_i, \alpha_i, \epsilon_i, p^{(i+1)}, \lambda, 0)$, $i = 1, \dots, k$, offer the same alternatives, but sequentially rather than in parallel.

2 Example. Consider the string transduction

$$\{ (a^{i_1} b a^{i_2} b \dots a^{i_n} b a^{i_{n+1}}, a^{i_1} b^{i_1} a^{i_2} b^{i_2} \dots a^{i_n} b^{i_n} a^{i_{n+1}}) \mid n \geq 0, i_1, \dots, i_{n+1} \geq 0 \}.$$

An obvious deterministic 2gsm reads each segment of a 's from left to right while copying it to the output. When encountering a b it rereads the segment from right to left. This second pass it writes b 's to the output tape.

This machine can be implemented by taking $\Sigma_1 = \Sigma_2 = \{a, b\}$, $Q = \{0, 1, 2, 3, 4, 5\}$, $q_{in} = 0$, $q_f = 5$, and δ consisting of the instructions

$$\begin{aligned} (0, \vdash, 1, \lambda, +1, 0, \lambda, 0) \\ (1, a, 1, a, +1, 2, \lambda, 0) \\ (2, b, 3, \lambda, -1, 5, \lambda, 0) \\ (3, a, 3, b, -1, 4, \lambda, +1) \\ (4, a, 4, \lambda, +1, 1, \lambda, +1) \end{aligned}$$

Note that the last three elements of the first instruction are irrelevant.

The computation of the 2gsm on input $aaabbaba$ can be visualized as in Figure 1, where we have labelled the edges of the computation by the strings that are written to the output (with λ omitted, for convenience). \square

Look-around. A *2gsm with regular look-around* (2gsm-rla) extends the basic 2gsm model, by allowing more complicated tests. In an instruction $(p, t, q_1, \alpha_1, \epsilon_1, q_0, \alpha_0, \epsilon_0) \in \delta$ all components are as before for the 2gsm, except the test t , which does not consist of a single letter σ , but of a triple

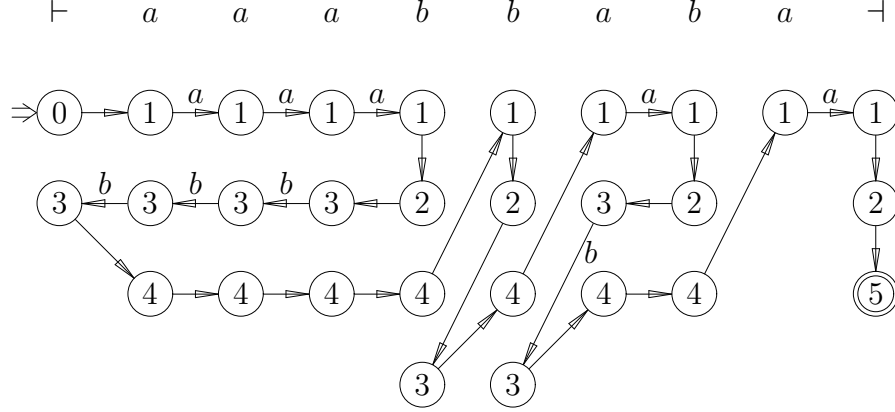


Figure 1: Computation for (a^3b^2aba, a^3b^3aba) of 2gsm from Example 2

$t = (R_\ell, \sigma, R_r)$, where $\sigma \in (\Sigma_1 \cup \{\vdash, \dashv\})$, and R_ℓ, R_r are regular languages such that $R_\ell, R_r \subseteq (\Sigma_1 \cup \{\vdash, \dashv\})^*$. This test t is satisfied if σ is the symbol under the reading head, and the strings to the left and the right of the head belong to R_ℓ and R_r respectively.

Obviously, it suffices to have tests (R_ℓ, σ, R_r) such that $R_\ell \cdot \sigma \cdot R_r \subseteq \vdash \Sigma_1^* \dashv$. For a given 2gsm-rla, an equivalent 2gsm-rla with that property is obtained by changing each test (R_ℓ, σ, R_r) into (R'_ℓ, σ, R'_r) where $R'_\ell = R_\ell \cap \vdash \Sigma_1^*$ (with the exception that $R'_\ell = \{\lambda\}$ when $\sigma = \vdash$), and similarly for R'_r . We observe here that this notion of ‘regular look-around’ generalizes the well-known notion of regular look-ahead for one-way automata (see, e.g., [Nij82, Eng77]).

Mso instructions. For a 2gsm with *mso-instructions* (2gsm-mso) the test and the moves of each instruction are given by mso formulas. To be precise, for $(p, t, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0) \in \delta$, t is given as a formula $\varphi(x)$ in $\text{MSO}(\Sigma_1 \cup \{\vdash, \dashv\}, *)$ with one free node variable x , and the moves μ_i are given by functional formulas $\varphi_i(x, y)$ in $\text{MSO}(\Sigma_1 \cup \{\vdash, \dashv\}, *)$ with two free node variables x and y (see below for the meaning of ‘functional’).

A test $t = \varphi(x)$ is evaluated for the string on the input tape with x valuated as the position taken by the reading head; more precisely, as our logic is defined for graphs, t is true whenever $\text{nd-gr}(\vdash w \dashv) \models \varphi(u)$, where w is the input string, and u is the node corresponding to the position of the reading head.

The 2gsm-mso does not move step-wise on the input tape, but it ‘jumps’ as specified by the formulas $\varphi_i(x, y)$, as follows. Assuming the machine is in position u , it moves to a position v for which $\text{nd-gr}(\vdash w \dashv) \models \varphi_i(u, v)$, where we have identified positions on the input tape with their corresponding nodes of the graph $\text{nd-gr}(\vdash w \dashv)$.

To guarantee that the $\varphi_i(x, y)$ describe deterministic moves of the reading head, we require that the relations specified by $\varphi_i(x, y)$ are functional, for each input string w , i.e., for every position u there is at most one position v such that $\text{nd-gr}(\vdash w \dashv) \models \varphi_i(u, v)$. Note that functionality is expressible in the logic: $(\forall x)(\forall y_1)(\forall y_2)[\varphi_i(x, y_1) \wedge \varphi_i(x, y_2) \rightarrow y_1 = y_2]$. Consequently, it is decidable; we may use Büchi’s result (Proposition 1, which is effective) to verify that it is satisfied by every string in $\vdash \Sigma_1^* \dashv$.

3 Example. Consider again the string transduction $m =$

$$\{ (a^{i_1} b a^{i_2} b \dots a^{i_n} b a^{i_{n+1}}, a^{i_1} b^{i_1} a^{i_2} b^{i_2} \dots a^{i_n} b^{i_n} a^{i_{n+1}}) \mid n \geq 0, i_1, \dots, i_{n+1} \geq 0 \}.$$

We use the predicate $\text{next}_a(x, y)$ to specify the first position y following x that is labelled by a :

$$x \prec y \wedge \text{lab}_a(y) \wedge (\forall z) [(x \prec z \wedge z \prec y) \rightarrow \neg \text{lab}_a(z)]$$

Similarly we construct an expression $\text{fis}_a(x, y)$ denoting the first a in the present segment of a ’s,

$$y \preceq x \wedge (\forall z)(y \preceq z \wedge z \preceq x \rightarrow \text{lab}_a(z)) \wedge \neg(\exists z)(\text{edge}_*(z, y) \wedge \text{lab}_a(z))$$

Using these predicates we build a deterministic 2gsm-mso that realizes m . In state 1 it walks along a segment of a ’s, copying it to the output tape. Then, when the segment is followed by a b , it jumps back to the first a of the segment for a second pass, in state 2. When the end of the segment is reached for the second time, the machine jumps to the next segment, returning to state 1. At the last a of the input the machine jumps to the right end marker, and halts in the final state 3.

Let $\Sigma_1 = \Sigma_2 = \{a, b\}$, $Q = \{1, 1', 2, 2', 3\}$, $q_{in} = 2'$, $q_f = 3$, and δ consisting of the transitions

$$\begin{aligned} & (1, (\exists y)(\text{edge}_*(x, y) \wedge \text{lab}_a(y)), 1, a, \text{edge}_*(x, y), 1', \lambda, x = y) \\ & (1', (\exists y)(\text{edge}_*(x, y) \wedge \text{lab}_b(y)), 2, b, \text{fis}_a(x, y), 3, \lambda, \text{lab}_\dashv(y)) \end{aligned}$$

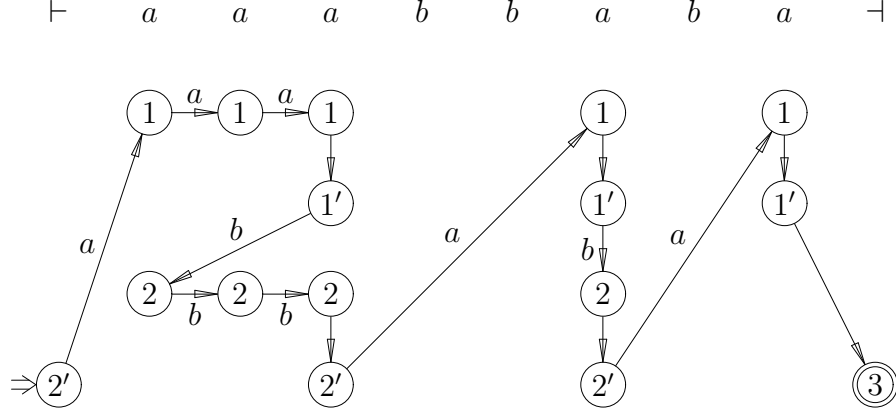


Figure 2: Computation for (a^3b^2aba, a^3b^3aba) of 2gsm-mso from Example 3

$$\begin{aligned}
 &(2, (\exists y)(\text{edge}_*(x, y) \wedge \text{lab}_a(y)), 2, b, \text{edge}_*(x, y), 2', \lambda, x = y) \\
 &(2', (\exists y)(x \prec y \wedge \text{lab}_a(y)), 1, a, \text{next}_a(x, y), 3, \lambda, \text{lab}_{\neg}(y))
 \end{aligned}$$

The computation of the machine on input a^3b^2aba can be visualized as in Figure 2 (where, again, λ is omitted from the edges of the computation). \square

Without loss of generality we assume that the 2m's we consider never write more than one symbol at a time, i.e., for each instruction $(p, \sigma, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0)$ we have $|\alpha_i| \leq 1$ (for $i = 1, 0$).

We abbreviate deterministic 2m's by adding a 'd' to the usual abbreviation, hence we speak of 2dgsm, 2dgsm-rla, and 2dgsm-mso. The families of string transductions realized by these three types of deterministic sequential machines are denoted by 2DGSM, 2DGSM^{RLA}, and 2DGSM^{MSO}, respectively.

Unlike their nondeterministic counterparts ([Kie75], see also Lemma 26 and the remark following it), deterministic 2gsm's are closed under composition, as was demonstrated by Chytil and Jákł. As an essential part of the proof the fact is used (proved in [HoU167]) that a 2dgsm can keep track of the state of another (deterministic) one-way finite state automaton working on the same tape (from left to right or from right to left). For the left-to-right case, it is clear how to do this as long as the reading head moves to the right. Backtracking ('undoing' a move) on the occasion of a step to the left, needs a rather ingenious back and forth simulation of the automaton.

4 Proposition ([ChJá77]) *2DGSM is closed under composition.*

In the remainder of this section we show that the three types of deterministic machines defined above are all equivalent, i.e., that $2DGSM = 2DGSM^{RLA} = 2DGSM^{MSO}$.

Every 2gsm is of course a simple 2gsm-rla, using trivial look-around tests, i.e., tests of the form (R_ℓ, σ, R_r) , with $R_\ell = \vdash \Sigma_1^*$, and $R_r = \Sigma_1^* \dashv$ (with the exceptions $R_\ell = \{\lambda\}$ when $\sigma = \vdash$, and $R_r = \{\lambda\}$ when $\sigma = \dashv$).

It follows from Büchi's result, Proposition 1, that any 2gsm-rla can be reinterpreted as a 2gsm-mso by changing the specification of the tests and moves into formulas, as follows.

First, consider a look-around test $t = (R_\ell, \sigma, R_r)$. Let $\psi_\ell(x)$ be a formula expressing that the string to the left of position x belongs to the regular language R_ℓ . It can be obtained from a closed formula ψ defining R_ℓ by restricting quantification to the positions to the left of x , i.e., by replacing subformulas $(\exists y)\xi(y)$ by $(\exists y)(y \prec x \wedge \xi(y))$ and $(\exists Y)\xi(Y)$ by $(\exists Y)((\forall y)(y \in Y \rightarrow y \prec x) \wedge \xi(Y))$.

Similarly, we obtain a formula $\psi_r(x)$ expressing that the string to the right of position x belongs to the regular language R_r . Clearly, the test t is equivalent to the formula $\varphi_t(x) = \psi_\ell(x) \wedge \text{lab}_\sigma(x) \wedge \psi_r(x)$.

Finally, one-step moves are easily translated into formulas. A move $\epsilon = +1$ is equivalent to stating that the new position is next to the original: $\text{edge}_*(x, y)$. Of course, $\epsilon = -1$ is symmetric, whereas $\epsilon = 0$ is expressed by $x = y$. Note that these formulas are functional.

These observations prove the first relations between the families of transductions.

5 Lemma. $2DGSM \subseteq 2DGSM^{RLA} \subseteq 2DGSM^{MSO}$.

The feature of 2dgsms that they can keep track of the state of a one-way finite state automaton (cf. the remark before Proposition 4), is modelled by us as regular look-around. Thus, for readers familiar with this feature it should be quite obvious that $2DGSM^{RLA} \subseteq 2DGSM$. Here we prove it using Proposition 4.

6 Lemma. $2DGSM^{RLA} \subseteq 2DGSM$.

Proof. By Proposition 4, 2DGSM is closed under composition. We prove the lemma by decomposing a given 2dgsm-rla \mathcal{M} into a series of 2dgsm's, together realizing the transduction of \mathcal{M} .

The final 2dgsm performs the required transduction, whereas all the other transductions ‘preprocess the tape’, by adding to the original input the outcome of the various tests of \mathcal{M} . As we also need this information for the positions containing the end-of-tape markers \vdash and \dashv , we start by a transduction that maps input w to the string $\triangleright w \triangleleft$, where \triangleright and \triangleleft are new symbols. Information concerning the end-of-tape positions is added to these new symbols. The other machines may ignore \vdash and \dashv , and treat \triangleright and \triangleleft as if they were these end-of-tape markers.

For each look-around test $t = (R_\ell, \sigma, R_r)$ of \mathcal{M} we introduce a 2dgsm \mathcal{M}_t that copies the input, while adding to each position the outcome of the test t for that position in the original string (ignoring any other additional information a previous transduction added to the string). The machine \mathcal{M}_t itself can be seen as the work of three consecutive 2dgsm's. The first one, simulating a finite state automaton recognizing R_ℓ , checks on each position whether the prefix read belongs to R_ℓ . It adds this information to the symbol at that position. The second transducer, processing the input from right to left, simulating a finite state automaton for the mirror image of R_r , adds information concerning the suffix. Note that the input has been reversed in the process. This can be undone by another reversal performed by a third 2dgsm.

Once the value of each look-around test of \mathcal{M} is added to the original input string, obviously the transduction of \mathcal{M} can be simulated by an ordinary 2dgsm. \square

Büchi's result (Proposition 1) allows us to show that the 2gsm-mso can be simulated by the 2gsm-rla. Additionally we need the following (folklore) result on the structure of certain regular languages (cf. [Pix96, Lemma 8.1]).

7 Lemma. *Let $\Delta \subseteq \Sigma$ be alphabets, and let $R \subseteq \Sigma^*$ be a regular language such that each string of R contains exactly one occurrence of a symbol from Δ . Then we may write R as a finite union of disjoint languages $R_\ell \cdot a \cdot R_r$, where $a \in \Delta$, and $R_\ell, R_r \subseteq (\Sigma - \Delta)^*$ are regular languages.*

Proof. Let \mathcal{A} be a deterministic finite automaton accepting R . Every path (in the state transition diagram of \mathcal{A}) from the initial state to a final state

passes exactly one transition labelled by a symbol from Δ . For any such transition (p, a, q) of \mathcal{A} let R_ℓ consist of all strings that label a path starting in the initial state of \mathcal{A} and ending in p , and symmetrically, let R_r consist of all strings that label a path from q to one of the final states of \mathcal{A} . Obviously, R_ℓ and R_r are regular, and R is the union of the languages $R_\ell \cdot a \cdot R_r$ taken over all such transitions. Since \mathcal{A} is deterministic, these languages are easily seen to be disjoint. \square

8 Lemma. $2\text{DGSM}^{\text{MSO}} \subseteq 2\text{DGSM}^{\text{RLA}}$.

Proof. We show how to simulate the instructions of a 2gsm-mso by a 2gsm-rla. Recall that such an instruction is specified as $(p, t, q_1, \alpha_1, \mu_1, q_0, \alpha_0, \mu_0)$, where t is a formula $\varphi(x)$ with one free node variable, and the moves μ_i are (functional) formulas $\varphi_i(x, y)$ with two free node variables.

Tests: unary node predicates. Consider a test $\varphi(x)$ in $\text{MSO}(\Sigma_1 \cup \{\vdash, \dashv\}, *)$. It can easily be simulated by regular look-around tests. Identifying $(\Sigma_1 \cup \{\vdash, \dashv\}) \times \{0, 1\}^{\{x\}}$ with $(\Sigma_1 \cup \{\vdash, \dashv\}) \times \{0, 1\}$, consider the language $L(\varphi)$, which is regular by Proposition 1. As each string of this language contains exactly one symbol with 1 as its second component, it can be written as a finite union of languages $R_\ell \cdot (\sigma, 1) \cdot R_r$, with regular languages $R_\ell, R_r \subseteq ((\Sigma_1 \cup \{\vdash, \dashv\}) \times \{0\})^*$, and $\sigma \in \Sigma_1 \cup \{\vdash, \dashv\}$, see Lemma 7. This implies that the test $\varphi(x)$ can be simulated by a finite disjunction of the look-around tests (R'_ℓ, σ, R'_r) , where each R'_ℓ, R'_r is obtained from the corresponding R_ℓ, R_r by dropping the second component (the 0-part) of the symbols. Of course, this disjunction is computed by testing each of its alternatives consecutively.

Moves: binary node predicates. Once the test of an instruction is evaluated, one of its moves is executed, and the output is written. This move is given as a formula $\varphi(x, y)$, specifying a functional relation between the present position x and the next position y on the input. Where the 2dgsm-mso may ‘jump’ to its next position, independent of the relative positions of x and y , a 2dgsm-rla can only step to one of the neighbouring positions of the tape, and has to ‘walk’ to the next position when simulating this jump.

Before starting the excursion from x to y the 2dgsm-rla determines the direction (left, right, or stay) by evaluating the tests $(\exists y)(y \prec x \wedge \varphi(x, y))$, $(\exists y)(x \prec y \wedge \varphi(x, y))$, and $(\exists y)(x = y \wedge \varphi(x, y))$ using the method that we have explained above. Since $\varphi(x, y)$ is functional, at most one of these tests is true.

In the sequel we assume that our target position y lies to the left of the present position x , i.e., test $(\exists y)(y \prec x \wedge \varphi(x, y))$ is true. The right-case can be treated in an analogous way; the stay-case is trivial.

Similarly to the case of tests, identify $(\Sigma_1 \cup \{\vdash, \dashv\}) \times \{0, 1\}^{\{x, y\}}$ with $(\Sigma_1 \cup \{\vdash, \dashv\}) \times \{0, 1\}^2$, and consider $L(y \prec x \wedge \varphi(x, y))$. Each string of this language contains exactly one symbol with $(0, 1)$ as its second component, the position of y , and it precedes a unique symbol with $(1, 0)$ as its second component, the position of x ; all other symbols carry $(0, 0)$. It can be written as a finite disjoint union of languages $R_\ell \cdot (\sigma, 0, 1) \cdot R_m \cdot (\tau, 1, 0) \cdot R_r$, with regular languages $R_\ell, R_m, R_r \subseteq ((\Sigma_1 \cup \{\vdash, \dashv\}) \times \{(0, 0)\})^*$ and $\sigma, \tau \in \Sigma_1 \cup \{\vdash, \dashv\}$, by applying Lemma 7 twice.

Our moves are functional, meaning that there is a unique position y that satisfies the predicate $\varphi(x, y)$ with x the present position. Still before starting the excursion from x to the new position y , the 2dgsm-rla determines which language in the union above describes this position by performing the regular look-around tests $(R'_\ell \cdot \sigma \cdot R'_m, \tau, R'_r)$, where each R'_ℓ, R'_m, R'_r is obtained from the corresponding R_ℓ, R_m, R_r by deleting the second component (the $(0, 0)$ -part) of the symbols.

The 2dgsm-rla now moves to the left. In each step it checks whether the segment of the input string between the present position (candidate y) and the starting position (corresponding to x) belongs to the regular language R'_m . This can be done by simulating a finite automaton for (the mirror image of) R'_m in the finite state control.

Each time this segment belongs to R'_m , it performs the rla-test $(R'_\ell, \sigma, \Sigma_1^* \dashv)$, to verify the requirement on the initial segment of the input. Once this last test is satisfied, it has found the position y and writes the output string. \square

We summarize.

9 Theorem. $2\text{DGSM} = 2\text{DGSM}^{\text{RLA}} = 2\text{DGSM}^{\text{MSO}}$.

A similar result can be obtained for nondeterministic gsm's by the same line of reasoning. However, in Lemma 6 we need the inclusion $2\text{DGSM} \circ 2\text{NGSM} \subseteq 2\text{NGSM}$ rather than $2\text{DGSM} \circ 2\text{DGSM} \subseteq 2\text{DGSM}$ (Proposition 4). This new inclusion can be proved like the latter one [ChJá77].

3 MSO Definable String Transductions

As explained in the Preliminaries, we consider mso logic on graphs as a means of specifying string transductions, rather than dealing directly with strings.

Although we are mainly interested in graph transductions that have string-like graphs as their domain and range, occasionally we find it useful to allow more general graphs as intermediate products of our constructions.

In this section we recall the definition of mso graph transductions, and from it we derive two families of mso definable string transductions, which differ in the way strings are represented by graphs. We present basic examples, and characterize the relation between the two families we have defined.

We start with the general definition.

An *mso definable transduction* [Cou91, Cou94, Eng91a, EnOo97, See92] is a (partial) function that constructs for a given input graph a new output graph as specified by a number of mso formulas. Here we consider the deterministic (or, ‘parameterless’) mso transductions of [Cou94]. For a graph satisfying a given domain formula φ_{dom} we take copies of each of the nodes, one for each element of a finite copy set C . The label of the c -copy of node x ($c \in C$) is determined by a set of formulas $\varphi_{\sigma}^c(x)$, one for each symbol σ in the output alphabet. We keep only those copies of the nodes for which exactly one of the label formulas is true. Edges are defined according to formulas $\varphi_{\gamma}^{c_1, c_2}(x, y)$: we construct an edge with label γ in the output graph from the c_1 -copy of x to the c_2 -copy of y whenever such a formula holds.

10 Definition. An *mso definable (graph) transduction* $\tau : \text{GR}(\Sigma_1, \Gamma_1) \rightarrow \text{GR}(\Sigma_2, \Gamma_2)$ is specified by

- a closed *domain formula* φ_{dom} ,
- a finite *copy set* C ,
- *node formulas* $\varphi_{\sigma}^c(x)$, with one free node variable x , for every $\sigma \in \Sigma_2$ and every $c \in C$, and
- *edge formulas* $\varphi_{\gamma}^{c_1, c_2}(x, y)$ with two free node variables x, y , for every $\gamma \in \Gamma_2$ and all $c_1, c_2 \in C$,

where all formulas are in $\text{MSO}(\Sigma_1, \Gamma_1)$.

For $g \in \text{GL}(\varphi_{\text{dom}})$ with node set V_g , the image $\tau(g)$ is the graph (V, E, ℓ) , defined as follows. We will write u^c rather than (u, c) for elements of $V_g \times C$.

- $V = \{u^c \mid u \in V_g, c \in C, \text{ there is exactly one } \sigma \in \Sigma_2 \text{ such that } g \models \varphi_\sigma^c(u)\},$
- $E = \{(u^{c_1}, \gamma, v^{c_2}) \mid u^{c_1}, v^{c_2} \in V, \gamma \in \Gamma_2, g \models \varphi_\gamma^{c_1, c_2}(u, v)\}, \text{ and}$
- $\ell(u^c) = \sigma \text{ if } g \models \varphi_\sigma^c(u), \text{ for } u^c \in V, \sigma \in \Sigma_2.$

□

11 Example. Let $\Sigma = \{a, b\}$. As a simple example we present an mso graph transduction from $\text{GR}(\Sigma, *)$ to $\text{GR}(*, \{a, b, *\})$ that transforms a linear graph representing a string into a ladder, while moving the symbols from the nodes to the steps.

Domain formula φ_{dom} expresses that the input graph is a string representation (see the end of Section 1).

The copy set C is $\{1, 2\}$.

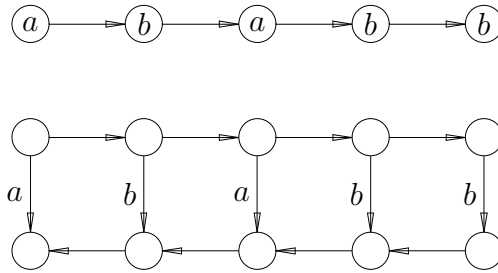
Each node is copied twice: $\varphi_*^1 = \varphi_*^2 = \text{true}$.

Unlabelled edges are copied twice, one of these in reverse:

$\varphi_*^{1,1} = \text{edge}_*(x, y), \varphi_*^{2,2} = \text{edge}_*(y, x), \varphi_*^{1,2} = \varphi_*^{2,1} = \text{false}.$

Labelled edges are introduced:

$\varphi_\sigma^{1,2} = (x = y) \wedge \text{lab}_\sigma(x), \varphi_\sigma^{1,1} = \varphi_\sigma^{2,1} = \varphi_\sigma^{2,2} = \text{false}, \text{ for } \sigma = a, b.$



□

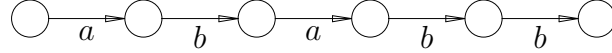
The family of mso definable graph transductions is denoted by grMSO . Its basic properties are summarized below, see, e.g., [Cou97, Prop. 5.5.6].

12 Proposition.

1. *grMSO is closed under composition.*
2. *The mso definable graph languages are closed under inverse mso definable graph transductions.*

We now consider mso definable graph transductions as a tool to specify string transductions.

There are two equally natural (and well-known) ways of representing a string as a graph. First, as we have seen in the Preliminaries, for a string $w \in \Sigma^*$ of length k , we may represent w by the graph $\text{nd-gr}(w)$ in $\text{GR}(\Sigma, *)$, consisting of k nodes labelled by the consecutive symbols of w , and $k - 1$ (unlabelled) edges representing the successor relation for the positions of the string. Dually, w can be represented by the graph $\text{ed-gr}(w)$ in $\text{GR}(*, \Sigma)$, consisting of $k + 1$ (unlabelled) nodes, connected by k edges that form a path labelled by the symbols of w . In the figure below we show $\text{ed-gr}(ababb)$. Note that $\text{ed-gr}(\lambda)$ consists of one unlabelled node.



It will turn out that the ‘edge graph representation’ of strings is more naturally related to two-way machines than the ‘node graph representation’.

13 Definition.

1. Let Σ_1, Σ_2 be two alphabets, and let $m \subseteq \Sigma_1^* \times \Sigma_2^*$ be a string transduction.
 - i. Its translation to graphs $\{(\text{ed-gr}(w), \text{ed-gr}(z)) \mid (w, z) \in m\}$ in $\text{GR}(*, \Sigma_1) \times \text{GR}(*, \Sigma_2)$ is denoted by $\text{ed-gr}(m)$;
 - ii. its translation to graphs $\{(\text{nd-gr}(w), \text{nd-gr}(z)) \mid (w, z) \in m\}$ in $\text{GR}(\Sigma_1, *) \times \text{GR}(\Sigma_2, *)$ is denoted by $\text{nd-gr}(m)$.
2. MSOS denotes the family of all string transductions m such that $\text{ed-gr}(m)$ belongs to grMSO, and MSOS_{nd} denotes the family of all string transductions m such that $\text{nd-gr}(m)$ belongs to grMSO.

□

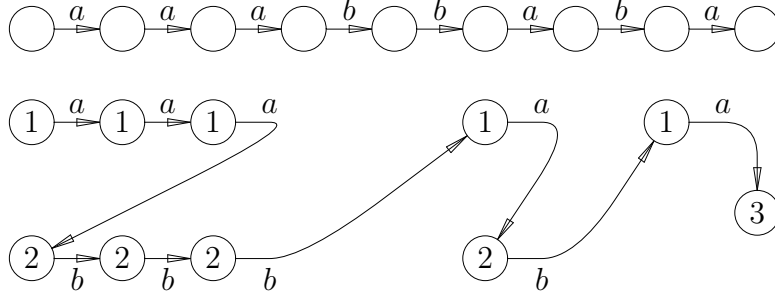


Figure 3: Edge representation for (a^3b^2aba, a^3b^3aba) , cf. Example 14

A transduction in MSOS is called an *mso definable string transduction*, and a transduction in MSOS_{nd} is called a λ -*restricted mso definable string transduction*. The reason for this terminology will be explained in Lemma 18.

14 Example. Consider the transduction $\text{ed-gr}(m)$, where m is the string transduction from Example 2,

$$\{ (a^{i_1}ba^{i_2}b \dots a^{i_n}ba^{i_{n+1}}, a^{i_1}b^{i_1}a^{i_2}b^{i_2} \dots a^{i_n}b^{i_n}a^{i_{n+1}}) \mid n \geq 0, i_1, \dots, i_{n+1} \geq 0 \}.$$

The formulas for the construction of the output graph have nodes as their reference points, whereas the information (symbols) is attached to the edges. Hence we frequently use the formula $\text{out}_\sigma(x) = (\exists y)\text{edge}_\sigma(x, y)$.

As in Example 3 we have an expression $\text{fis}'_a(x, y)$ denoting the first node in the present segment of a 's, this time referring to outgoing edges:

$$y \preceq x \wedge (\forall z)(y \preceq z \wedge z \preceq x \rightarrow \text{out}_a(z)) \wedge \neg(\exists z)(\text{edge}_a(z, y))$$

Similarly, we have the edge variant $\text{next}'_a(x, y)$ by replacing the subformulas $\text{lab}_a(y)$ by $\text{out}_a(y)$ in the original formula $\text{next}_a(x, y)$.

Choosing the copy set $C = \{1, 2, 3\}$, and the domain formula defining edge representations of strings, the transduction $\text{ed-gr}(m)$ is defined by the following formulas.

$$\begin{aligned} \varphi_*^1 &= \text{out}_a(x) \\ \varphi_*^2 &= \text{out}_a(x) \wedge (\exists y)(x \preceq y \wedge \text{out}_b(y)) \\ \varphi_*^3 &= \neg \text{out}_a(x) \wedge \neg \text{out}_b(x), \quad \text{the final node of the string,} \\ \varphi_a^{1,1} &= \text{edge}_a(x, y) \end{aligned}$$

$$\begin{aligned}
\varphi_a^{1,2} &= (\exists z)(\text{edge}_a(x, z) \wedge \neg \text{out}_a(z)) \wedge \text{fis}'_a(x, y) \\
\varphi_a^{1,3} &= \neg(\exists z)(\varphi_a^{1,1}(x, z) \vee \varphi_a^{1,2}(x, z)) \\
\varphi_b^{2,2} &= \text{edge}_a(x, y) \\
\varphi_b^{2,1} &= (\exists z)(\text{edge}_a(x, z) \wedge \neg \text{out}_a(z)) \wedge \text{next}'_a(x, y) \\
\varphi_b^{2,3} &= \neg(\exists z)(\varphi_b^{2,1}(x, z) \vee \varphi_b^{2,2}(x, z)) \\
\varphi_\sigma^{3,j} &= \text{false, for } j = 1, 2, 3.
\end{aligned}$$

The construction is illustrated in Figure 3 for $(a^3b^2aba, a^3b^3aba) \in m$. Note that we have put the copy numbers within the nodes. \square

The transition from one graph representation to the other is (essentially) definable as mso graph transduction, and will be heavily used in the sequel. We discuss this in the next example.

15 Example. The graph transduction $\text{ed2nd} = \{ (\text{ed-gr}(w), \text{nd-gr}(w)) \mid w \in \Sigma^* \} : \text{GR}(*, \Sigma) \rightarrow \text{GR}(\Sigma, *)$ from the edge representation of a string into its node representation is mso definable, as follows.

- φ_{dom} expresses that the input is a string representation, an edge-labelled path (consisting of at least one node);
- the copy set C equals $\{1\}$;
- $\varphi_\sigma^1 = (\exists y)(\text{edge}_\sigma(x, y))$, i.e., the label σ is moved from the edge to its source node. None of these formulas is true for the final node of the input graph, which means that this node is not copied;
- $\varphi_*^{1,1} = \bigvee_{\sigma \in \Sigma} \text{edge}_\sigma(x, y)$, i.e., edges are copied, without their labels.

The inverse mapping $\text{ed2nd}^{-1} = \{ (\text{nd-gr}(w), \text{ed-gr}(w)) \mid w \in \Sigma^* \} : \text{GR}(\Sigma, *) \rightarrow \text{GR}(*, \Sigma)$ is *not* mso definable: The representation $\text{nd-gr}(\lambda)$ of the empty string has no nodes that can be copied to obtain the single node of $\text{ed-gr}(\lambda)$.

If we omit the empty string, the graph transduction $\text{nd2ed} = \{ (\text{nd-gr}(w), \text{ed-gr}(w)) \mid w \in \Sigma^*, w \neq \lambda \}$ can be defined as follows.

- φ_{dom} again expresses that the input is a string representation, a (non-empty) node-labelled path;

- the copy set equals $\{1, 2\}$;
- $\varphi_*^1 = \text{true}$, $\varphi_*^2 = \neg(\exists y)(\text{edge}_*(x, y))$, i.e., all nodes are copied once, except the last one which gets two copies;
- $\varphi_\sigma^{1,1} = \text{edge}_*(x, y) \wedge \text{lab}_\sigma(x)$, i.e., the label is moved from the node to its outgoing edge;
- $\varphi_\sigma^{1,2} = (x = y) \wedge \text{lab}_\sigma(x)$, which deals with the last edge;
- $\varphi_\sigma^{2,1} = \varphi_\sigma^{2,2} = \text{false}$.

□

The above example illustrates an important technical point: every mso graph transduction maps the empty graph to itself (provided it belongs to the domain). This means that, when using the node-encoding nd-gr for strings, the empty string can only be mapped to itself. As we do not want to restrict ourselves to this kind of transductions, we have chosen to consider both variants of mso definable string transductions. Although $\text{nd-gr}(w)$ is a slightly more direct graph representation of the string w in terms of its positions and their successor relation, the advantage of $\text{ed-gr}(w)$ is that it is never empty and thus satisfies all the usual logical laws.

The transition from node representation to edge representation for strings does not influence the validity of Büchi’s result.

16 Proposition. *A string language $K \subseteq \Sigma^*$ is regular iff there is a closed formula φ of $\text{MSO}(*, \Sigma)$ such that $K = \{w \in \Sigma^* \mid \text{ed-gr}(w) \models \varphi\}$.*

Proof. Rather direct, using Büchi’s result (Proposition 1(2)) and Proposition 12(2). We consider one implication (from right to left) only.

Let the string language $K \subseteq \Sigma^*$ be defined by the closed formula φ of $\text{MSO}(*, \Sigma)$, as in the statement of the lemma (using the edge representation). We show that there exists a formula defining K using the node representation. Consider the mso definable graph transduction nd2ed mapping $\text{nd-gr}(w)$ to $\text{ed-gr}(w)$ for all non-empty $w \in \Sigma^*$, cf. Example 15. The graph language $\text{nd2ed}^{-1}(GL(\varphi)) = \{\text{nd-gr}(w) \mid w \in \Sigma^*, w \neq \lambda, \text{ed-gr}(w) \models \varphi\}$ is mso definable, say by an mso formula ψ of $\text{MSO}(\Sigma, *)$. It defines the string language $L(\psi) = \{w \in \Sigma^* \mid \text{nd-gr}(w) \models \psi\} = \{w \in \Sigma^* \mid \text{ed-gr}(w) \models \varphi, w \neq \lambda\} = K - \{\lambda\}$. If $\lambda \notin K$, then we are done; otherwise, consider $L(\psi \vee \neg(\exists x)\text{true})$.

□

The families MSOS_{nd} and MSOS are equal, up to a small technicality involving the empty string—a point already illustrated in Example 15, and in the proof of Proposition 16.

To prove this, we use the following basic fact (cf. [Cou94, Proposition 3.3]).

17 Lemma. *Let τ_1 and τ_2 be mso definable graph transductions from $\text{GR}(\Sigma_1, \Gamma_1)$ to $\text{GR}(\Sigma_2, \Gamma_2)$.*

If τ_1 and τ_2 have disjoint domains, then also $\tau_1 \cup \tau_2 \in \text{grMSO}$.

Proof. Consider τ_i fixed by the copy set C_i and formulas $\varphi_{\text{dom},i}$, $\varphi_{\sigma,i}^c$, and $\varphi_{\gamma,i}^{c_1,c_2}$. We may assume that C_1 and C_2 are disjoint.

The domain formula for the union is the disjunction $\varphi_{\text{dom},1} \vee \varphi_{\text{dom},2}$; its copy set is $C = C_1 \cup C_2$.

The node formulas and the edge formulas for both transductions are also taken together (by disjunction), but we ensure that they are applicable only for the appropriate input by changing $\varphi_{\sigma,i}^c$ to $\varphi_{\text{dom},i} \wedge \varphi_{\sigma,i}^c$, and similarly for the edge formulas. We add $\varphi_{\gamma}^{c_1,c_2} = \varphi_{\gamma}^{c_2,c_1} = \text{false}$ for $c_1 \in C_1$, $c_2 \in C_2$, $\gamma \in \Gamma_2$. \square

18 Lemma. *Let $m \subseteq \Sigma_1^* \times \Sigma_2^*$ be a string transduction. Then $m \in \text{MSOS}_{\text{nd}}$ iff $m \in \text{MSOS}$ and $(\lambda, z) \in m$ implies $z = \lambda$.*

Proof. (1) From left to right; assume $m \in \text{MSOS}_{\text{nd}}$, i.e., $\text{nd-gr}(m) \in \text{grMSO}$. We split m into the mappings $\hat{m} = \{(w, z) \in m \mid z \neq \lambda\}$, and $m_\lambda = \{(w, z) \in m \mid z = \lambda\}$.

As $\text{nd-gr}(m) \in \text{grMSO}$, also $\text{ed-gr}(\hat{m}) = \text{ed2nd} \circ \text{nd-gr}(m) \circ \text{nd2ed}$ is mso definable, by Proposition 12(1).

By Proposition 12(2), the domain of $\text{ed-gr}(m_\lambda)$ is mso definable as it is the inverse image of $\{\text{nd-gr}(\lambda)\}$ for the transduction $\text{ed2nd} \circ \text{nd-gr}(m)$. Now it is easily seen that $\text{ed-gr}(m_\lambda) \in \text{grMSO}$ using for φ_{dom} the formula defining the domain of $\text{ed-gr}(m_\lambda)$, $C = \{1\}$, $\varphi_*^1 = \neg(\exists y)\text{edge}(x, y)$, and $\varphi_{\gamma}^{1,1} = \text{false}$.

The union $\text{ed-gr}(m) = \text{ed-gr}(\hat{m}) \cup \text{ed-gr}(m_\lambda)$ is mso definable by Lemma 17. Hence, $m \in \text{MSOS}$. We have discussed already that the image of λ under m must be λ (provided λ belongs to the domain of m) as $\text{nd-gr}(\lambda)$ has no nodes to copy.

(2) From right to left; assume $m \in \text{MSOS}$, i.e., $\text{ed-gr}(m) \in \text{grMSO}$.

Then also $\text{nd-gr}(\hat{m}) = \text{nd2ed} \circ \text{ed-gr}(m) \circ \text{ed2nd}$ is mso definable, where $\hat{m} = m - \{(\lambda, \lambda)\}$.

We are ready when λ does not belong to the domain of m . Otherwise, as the transduction $\{(\text{nd-gr}(\lambda), \text{nd-gr}(\lambda))\}$, mapping the empty graph to itself, is easily seen to be mso definable, $\text{nd-gr}(m) \in \text{grMSO}$ follows by Lemma 17. \square

We finally observe that, from Proposition 12(1), it immediately follows that MSOS is closed under composition. Together with the closure under composition of 2DGSM (Proposition 4) this has been a strong indication for the equality of these two families, proved in the next section.

4 Logic and Machines

In this section we establish our main result, the equivalence of the deterministic two-way sequential machines from Section 2, and the mso definable string transductions from Section 3: $\text{MSOS} = 2\text{DGSM}$.

The first steps towards this result were taken already in Section 2 when we introduced the 2gsm with mso instructions, and showed its equivalence to the basic two-way generalized sequential machine.

One technical notion that will be essential to bridge the final gap between logic and machine is modelled after Figure 1 in Example 2. That figure depicts the computation of a 2gsm on a given input string. The input string w can naturally be represented by $\text{nd-gr}(\vdash w \dashv)$ with nodes corresponding to positions on the tape. On the other hand, the output string z is represented as $\text{ed-gr}(z')$ where the edges conveniently correspond to steps of the 2gsm from one position to another (and where z is obtained from z' by erasing λ , i.e., by removing the unlabelled edges).

We introduce a notation for this representation. Let $m : \Sigma_1^* \rightarrow \Sigma_2^*$ be a string transduction. We use $\text{tape}(m)$ to denote the graph transduction $\{ (\text{nd-gr}(\vdash w \dashv), \text{ed-gr}(z)) \mid (w, z) \in m \}$ from $\text{GR}(\Sigma_1 \cup \{\vdash, \dashv\}, *)$ to $\text{GR}(*, \Sigma_2)$.

19 Example. Consider the transduction $\text{tape}(m)$, where m is the string transduction from Example 2,

$$\{ (a^{i_1} b a^{i_2} b \dots a^{i_n} b a^{i_{n+1}}, a^{i_1} b^{i_1} a^{i_2} b^{i_2} \dots a^{i_n} b^{i_n} a^{i_{n+1}}) \mid n \geq 0, i_1, \dots, i_{n+1} \geq 0 \}.$$

Previously we have shown that $m \in 2\text{DGSM}$, here we will demonstrate that $\text{tape}(m)$ is an mso definable graph transduction.

Recall the predicate $\text{next}_\sigma(x, y)$ from Example 3.

For $\text{tape}(m)$ the domain formula specifies linear graphs of the form $\text{nd-gr}(\vdash w \dashv)$, $w \in \{a, b\}^*$, the copy set C is $\{1, 3, 5\}$, and we have formulas

$$\begin{aligned} \varphi_*^1 &= \text{lab}_a(x), \\ \varphi_*^3 &= \text{lab}_a(x) \wedge (\exists y)(x \preceq y \wedge \text{lab}_b(y)), \\ \varphi_*^5 &= \text{lab}_{\dashv}(x), \\ \varphi_a^{1,1} &= \text{edge}_*(x, y), \\ \varphi_a^{1,3} &= (x = y) \wedge \neg(\exists z)(\text{edge}_*(x, z) \wedge \text{lab}_a(z)), \\ \varphi_a^{1,5} &= \text{edge}_*(x, y), \end{aligned}$$

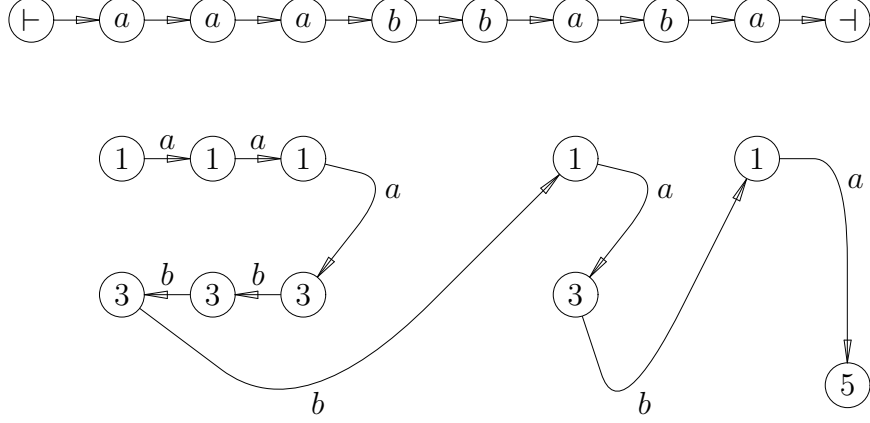


Figure 4: Mso transduction $\text{tape}(m)$ from Example 19

$$\begin{aligned}
\varphi_b^{3,3} &= \text{edge}_*(y, x), \\
\varphi_b^{3,1} &= (\exists z)(\text{next}_b(x, z) \wedge \text{next}_a(z, y)) \wedge \neg(\exists z)(\text{edge}_*(z, x) \wedge \text{lab}_a(z)), \\
&\text{i.e., connect to the first } a \text{ of the next segment when we are at the first} \\
&\text{ } a \text{ of the present segment,} \\
\varphi_b^{3,5} &= \neg(\exists z)(\varphi_b^{3,1}(x, z) \vee \varphi_b^{3,3}(x, z)), \\
\varphi_\sigma^{i,j} &= \text{false, in all other cases.}
\end{aligned}$$

Note that the output of the transduction (cf. the lower graph in Figure 4) is obtained by contracting unlabelled paths in the computation graph of the 2dgsM from Example 2, Figure 1. \square

The observation from the example is generally true: a string transduction m is realized by a 2dgsM if and only if its graph representation $\text{tape}(m)$ is mso definable. We prove the two implications separately.

20 Lemma. *Let $m : \Sigma_1^* \rightarrow \Sigma_2^*$ be a string transduction. If $m \in \text{2DGSM}$, then $\text{tape}(m) \in \text{grMSO}$.*

Proof. Let $\mathcal{M} = (Q, \Sigma_1, \Sigma_2, \delta, q_{in}, q_f)$ be a 2dgsM realizing the string transduction $m : \Sigma_1^* \rightarrow \Sigma_2^*$, and consider a fixed input string $w = \sigma_1 \cdots \sigma_n$, $\sigma_i \in \Sigma_1$ for $i = 1, \dots, n$. Additionally we use $\sigma_0 = \vdash$ and $\sigma_{n+1} = \dashv$.

We can visualize the ‘computation space’ of \mathcal{M} on w by constructing a graph $\gamma_{\mathcal{M}}(w)$ that has as its nodes the pairs $\langle p, i \rangle$, where p is a state of \mathcal{M} ,

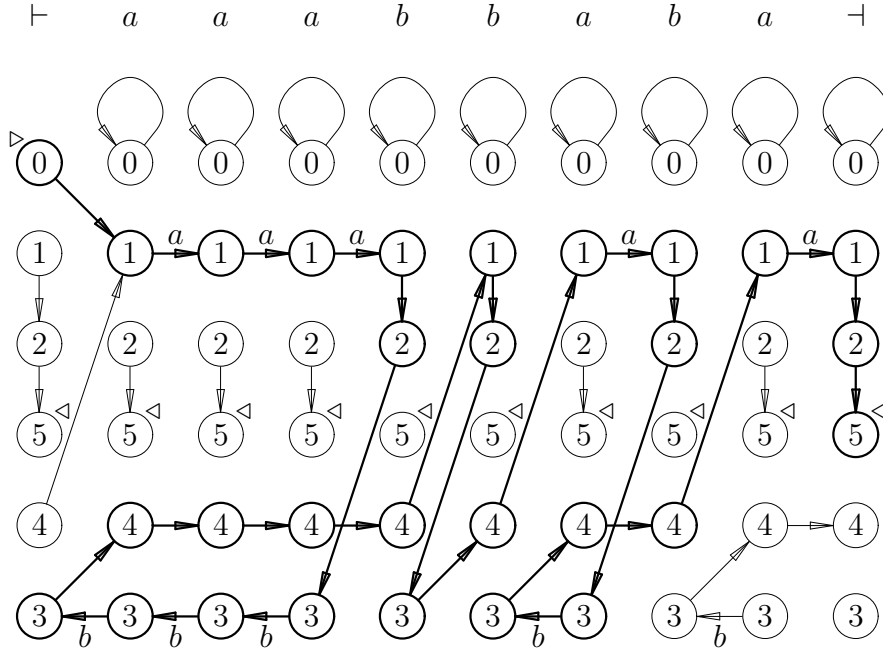


Figure 5: Computation space $\gamma_{\mathcal{M}}(a^3b^2aba)$ for the 2dgsM \mathcal{M} in Example 2

and $i \in \{0, 1, \dots, n, n+1\}$ is one of the positions of the input tape carrying $\vdash w \dashv$. The edges of $\gamma_{\mathcal{M}}(w)$ are chosen in accordance with the instruction set δ of \mathcal{M} : for each instruction $t = (p, \sigma, q_1, \alpha_1, \epsilon_1, q_0, \alpha_0, \epsilon_0)$ in δ there is an edge from $\langle p, i \rangle$ to $\langle q_1, i + \epsilon_1 \rangle$ if σ_i equals σ , and an edge from $\langle p, i \rangle$ to $\langle q_0, i + \epsilon_0 \rangle$ otherwise. The edge is labelled by the output symbol $\alpha_i \in \Sigma_2 \cup \{\lambda\}$. In this context we will consider λ as a labelling symbol (rather than as a string of length zero) in order to avoid notational complications.

In Figure 5 we illustrate the computation space for the 2dgsM from Example 2 on input a^3b^2aba (with output λ omitted, as usual). The computation on that input is represented as a bold path (cf. Figure 1).

As \mathcal{M} is deterministic, every node of $\gamma_{\mathcal{M}}(w)$ has at most one outgoing edge. The output of the computation of \mathcal{M} on w can then be read from $\gamma_{\mathcal{M}}(w)$ by starting in node $\langle q_{in}, 0 \rangle$, representing \mathcal{M} in its initial configuration, and following the path along the outgoing edges. The computation is successful if it ends in a final configuration $\langle q_f, k \rangle$. We will mark the initial and final nodes of $\gamma_{\mathcal{M}}(w)$ by special labels \triangleright and \triangleleft , the other nodes remain unlabelled (represented in our specification by ‘*’).

Note that the graph $\gamma_{\mathcal{M}}(w)$ does not only represent the computation of

\mathcal{M} on w starting in the initial state and 0-th position of the tape (marked by \vdash) but rather all possible computations that result from placing \mathcal{M} on an arbitrary position of the tape, in an arbitrary state.

We construct a series of mso graph transductions, the composition of which maps $\text{nd-gr}(\vdash w \dashv)$ to $\text{ed-gr}(z)$ for each $(w, z) \in m$. As grMSO is closed under composition (Proposition 12), this proves the lemma.

The first graph transduction τ_1 maps $\text{nd-gr}(\vdash w \dashv)$ to $\gamma_{\mathcal{M}}(w)$. The second graph transduction τ_2 selects the path in $\gamma_{\mathcal{M}}(w)$ corresponding to the successful computation of \mathcal{M} on w (if it exists) by keeping only those nodes that are reachable from the initial configuration and lead to a final configuration. The last graph transduction τ_3 removes edges labelled by λ (used as a symbol representing the empty string) while contracting paths consisting of these edges.

Step one: constructing $\gamma_{\mathcal{M}}(w)$. Let $\tau_1 : \text{GR}(\Sigma_1 \cup \{\vdash, \dashv\}, *) \rightarrow \text{GR}(\{*, \triangleright, \triangleleft\}, \Sigma_2 \cup \{\lambda\})$ be the graph transduction that constructs $\gamma_{\mathcal{M}}(w)$. We follow the general description above, and formalize τ_1 as mso transduction.

The domain formula of the transduction specifies that the graph is of the form $\text{nd-gr}(\vdash w \dashv)$ for some string w . The copy set equals $C = Q$, where Q is the set of states of \mathcal{M} . The node $\langle q, i \rangle$ of $\gamma_{\mathcal{M}}(w)$ is identified with u_i^q , the q -copy of the node u_i of $\text{nd-gr}(\vdash w \dashv)$ corresponding to the i -th position of the input tape, labelled with σ_i .

The labels of the edges are chosen according to the instructions of \mathcal{M} . For $\alpha \in \Sigma_2 \cup \{\lambda\}$, $p, q \in Q$, and $\epsilon \in \{-1, 0, +1\}$ let $\text{step}[\epsilon]_{\alpha}^{p,q}(x)$ be the following disjunction, where the unspecified ‘dots’ range over their respective components:

$$\bigvee_{(p, \sigma, q, \alpha, \epsilon, \dots) \in \delta} \text{lab}_{\sigma}(x) \vee \bigvee_{\substack{(p, \tau, \dots, q, \alpha, \epsilon) \in \delta \\ \tau \neq \sigma}} \text{lab}_{\sigma}(x)$$

Then,

$$\begin{aligned} \varphi_{\alpha}^{p,q} = & (\text{edge}_*(x, y) \wedge \text{step}[+1]_{\alpha}^{p,q}(x)) \\ & \vee (x = y \wedge \text{step}[0]_{\alpha}^{p,q}(x)) \\ & \vee (\text{edge}_*(y, x) \wedge \text{step}[-1]_{\alpha}^{p,q}(x)) \end{aligned}$$

All copies of the nodes are present, with special labels for initial and final nodes:

$$\varphi_{\triangleright}^q = \text{lab}_{\triangleright}(x), \text{ when } q = q_{in}, \text{ and } \varphi_{\triangleright}^q = \text{false}, \text{ otherwise.}$$

$$\varphi_{\triangleleft}^q = \text{true}, \text{ when } q = q_f, \text{ and } \varphi_{\triangleleft}^q = \text{false}, \text{ otherwise.}$$

$$\varphi_*^q = \neg\varphi_{\triangleright}^q(x) \wedge \neg\varphi_{\triangleleft}^q(x).$$

Note that we assume that $q_{in} \neq q_f$, in order to avoid that both $\varphi_{\triangleright}^q$ and $\varphi_{\triangleleft}^q$ are defined for the initial node. This is the case when \mathcal{M} accepts any input in its initial state without executing instructions. We satisfy the assumption by adding additional instructions to a new final state.

Step two: selecting the computation path. The transduction $\tau_2 : \text{GR}(\{*, \triangleright, \triangleleft\}, \Sigma_2 \cup \{\lambda\}) \rightarrow \text{GR}(*, \Sigma_2 \cup \{\lambda\})$ removes nodes that are not on the path from the node labelled by \triangleright to a node labelled by \triangleleft (if it exists). Nodes that are not on such a path do not correspond to the configurations that are part of the (successful) computation of \mathcal{M} on w . Note that if such a path exists, then it is unique.

Recall that the predicate \preceq specifies the existence of a path from x to y . By $x \preceq_{\lambda} y$ we restrict ourselves below to a path containing only edges with label λ .

Formally,

$$\varphi_{\text{dom}} = (\exists x)(\exists y)[\text{lab}_{\triangleright}(x) \wedge \text{lab}_{\triangleleft}(y) \wedge x \preceq y],$$

$$C = \{1\},$$

$$\varphi_*^1(x) = (\exists y)(\exists z)[\text{lab}_{\triangleright}(y) \wedge y \preceq x \wedge \text{lab}_{\triangleleft}(z) \wedge x \preceq z]$$

$$\text{and, for } \alpha \in \Sigma_2 \cup \{\lambda\}, \varphi_{\alpha}^{1,1}(x, y) = \text{edge}_{\alpha}(x, y).$$

Step three: contracting λ -paths. The last graph transduction of three, $\tau_3 : \text{GR}(*, \Sigma_2 \cup \{\lambda\}) \rightarrow \text{GR}(*, \Sigma_2)$ deletes all nodes that have an outgoing λ -labelled edge, and contracts each λ -path to its last node.

This can be specified with the trivial copy set $C = \{1\}$, node formula $\varphi_*^1 = \neg(\exists y)(\text{edge}_{\lambda}(x, y))$, and edge formulas $\varphi_{\alpha}^{1,1} = (\exists z)(\text{edge}_{\alpha}(x, z) \wedge z \preceq_{\lambda} y)$, for $\alpha \in \Sigma_2$. \square

Now that the 2dgsm has learned to understand the language of monadic second-order logic, cf. Theorem 9, the converse of the previous result has a rather straightforward proof.

21 Lemma. *Let $m : \Sigma_1^* \rightarrow \Sigma_2^*$ be a string transduction. If $\text{tape}(m) \in \text{grMSO}$, then $m \in \text{2DGSM}$.*

Proof. Starting with the mso transduction $\text{tape}(m) : \text{GR}(\Sigma_1 \cup \{\vdash, \dashv\}, *) \rightarrow \text{GR}(*, \Sigma_2)$ we build a 2dgsm-mso \mathcal{M} for m that closely follows the mso specification of $\text{tape}(m)$.

Assume $\text{tape}(m)$ is specified by domain formula φ_{dom} , copy set C , node formulas φ_*^c , $c \in C$, and edge formulas $\varphi_\sigma^{c_1, c_2}$, $c_1, c_2 \in C$, $\sigma \in \Sigma_2$. The state set of \mathcal{M} is (in principle) equal to the copy set C : when $\varphi_\sigma^{c_1, c_2}(u, v)$ is true for a pair u, v of nodes, then \mathcal{M} , visiting the position corresponding to u of the input tape in state c_1 , may move to the position corresponding to v changing to state c_2 , while writing σ to the output tape.

Note that, for each input graph g , $\text{tape}(m)(g)$ defines a graph representation of a string, hence at most one of these formulas defines an edge in a given position (node) and a given state (copy). However, in general the formula $\varphi_\sigma^{c_1, c_2}$ is only functional as far as graphs g satisfying the domain formula φ_{dom} are concerned, and for these graphs only when restricted to nodes for which the respective c_1 and c_2 copies are defined. Since our formal definition of 2dgsm-mso demands functional moves, we consider the formulas $\psi_\sigma^{c_1, c_2}(x, y) = \varphi_\sigma^{c_1, c_2}(x, y) \wedge \varphi_*^{c_1}(x) \wedge \varphi_*^{c_2}(y) \wedge \varphi_{\text{dom}}$.

The instructions of \mathcal{M} are of the form

$$(c_1, (\exists y)(\psi_\sigma^{c_1, c_2}(x, y)), c_2, \sigma, \psi_\sigma^{c_1, c_2}(x, y))$$

– but this is 5-tuple notation, and has to be replaced by 8-tuples where for a fixed state c_1 each of the alternatives $(c_2, \sigma) \in C \times \Sigma_2$ has to be tested consecutively, as explained in the paragraph about 2gsm in Section 2 (using additional states).

If none of the edge formulas gives a positive result, the present node has no successor, which indicates the last position of the output string. In that case, the series of consecutive tests ends up in the final state q_f .

Initially \mathcal{M} has to find the unique node of the output graph that has no incoming edges. We solve this by adding the new initial state q_{in} from which this node is found by testing all possibilities, but again in a consecutive

fashion, for $c_2 \in C$:

$$(q_{in}, (\exists y)[\varphi_*^{c_2}(y) \wedge \neg \text{incom}^{c_2}(y)], c_2, \lambda, \varphi_*^{c_2}(y) \wedge \neg \text{incom}^{c_2}(y))$$

where $\text{incom}^{c_2}(y)$ abbreviates $(\exists z) \bigvee_{c_1 \in C, \sigma \in \Sigma_2} (\psi_\sigma^{c_1, c_2}(z, y))$. \square

22 Lemma. *Let Σ be an alphabet. The transduction $\text{tape}(id) : \text{GR}(\Sigma \cup \{\vdash, \dashv\}, *) \rightarrow \text{GR}(*, \Sigma)$ mapping $\text{nd-gr}(\vdash w \dashv)$ to $\text{ed-gr}(w)$ is an element of grMSO , as is its inverse $\text{tape}(id)^{-1}$.*

Proof. The identity on Σ^* is easily performed by an 2dgsm. Hence $\text{tape}(id) \in \text{grMSO}$, by Lemma 20.

As for the inverse $\text{tape}(id)^{-1}$, note that mapping $\text{ed-gr}(w)$ to $\text{ed-gr}(\vdash w \dashv)$ is mso definable because $\text{ed-gr}(w)$ has at least one node, which may be copied to provide the additional nodes that are connected by edges labelled by \vdash and \dashv to the original graph. We now compose this mapping by ed2nd , which is mso definable by Example 15. \square

We complete the section by deriving the equivalence between the mso definable string transductions and the deterministic two-way finite state transductions, uniting logic and machines.

23 Theorem. $\text{MSOS} = \text{2DGSM}$.

Proof. By our previous lemma, the transduction $\text{tape}(id)$ from $\text{nd-gr}(\vdash w \dashv)$ to $\text{ed-gr}(w)$, for $w \in \Sigma_1^*$, is an element of grMSO , as is its inverse $\text{tape}(id)^{-1}$. By the equalities $\text{tape}(m) = \text{tape}(id) \circ \text{ed-gr}(m)$, and $\text{ed-gr}(m) = \text{tape}(id)^{-1} \circ \text{tape}(m)$, and the closure of grMSO under composition (Proposition 12), we have $m \in \text{MSOS}$ iff (by definition) $\text{ed-gr}(m) \in \text{grMSO}$ iff $\text{tape}(m) \in \text{grMSO}$.

The result now follows from Lemmas 20 and 21 demonstrating $\text{tape}(m) \in \text{grMSO}$ iff $m \in \text{2DGSM}$. \square

As an immediate consequence of this result and Lemma 18 we obtain the equivalence between the corresponding λ -restricted transductions.

We use $\text{2DGSM}\lambda$ to denote those relations m in 2DGSM that satisfy $(\lambda, z) \in m$ implies $z = \lambda$, cf. Lemma 18.

24 Corollary. $\text{MSOS}_{\text{nd}} = \text{2DGSM}\lambda$.

5 Nondeterminism

In this section we define the nondeterministic mso definable graph transductions, and their derived string relatives. We observe that nondeterministic mso transductions are related to the deterministic mso transductions via relabelling of the input.

A nondeterministic variant of mso definable transductions is considered in [Cou91, Cou94]. All the formulas of the deterministic version may now have additional free node-set variables X_1, \dots, X_k , called ‘parameters’, the same for each of the formulas. For each valuation of the parameters (by sets of nodes of the input graph) that satisfies the domain formula, the other formulas define the output graph as before. Hence each valuation may lead to a different output graph for the given input graph: nondeterminism.

More formally, a *nondeterministic mso definable (graph) transduction* $\tau \subseteq \text{GR}(\Sigma_1, \Gamma_1) \times \text{GR}(\Sigma_2, \Gamma_2)$ is specified by

- a set of *parameters* X_1, \dots, X_k , $k \geq 0$,
- a *domain formula* $\varphi_{\text{dom}}(X_1, \dots, X_k)$,
- a finite *copy set* C ,
- *node formulas* $\varphi_\sigma^c(x, X_1, \dots, X_k)$ for $\sigma \in \Sigma_2$, $c \in C$, and
- *edge formulas* $\varphi_\gamma^{c_1, c_2}(x, y, X_1, \dots, X_k)$ for $\gamma \in \Gamma_2$, $c_1, c_2 \in C$,

where all formulas are in $\text{MSO}(\Sigma_1, \Gamma_1)$.

Recall from Section 1 that an input graph together with a valuation of the parameters can be represented by a Ξ -valuated graph g which has node labels in $\Sigma_1 \times \{0, 1\}^\Xi$ (where $\Xi = \{X_1, \dots, X_k\}$) such that $g|_{\Sigma_1}$ is the input graph, and ν_g is the valuation. By definition, $g \in GL(\varphi_{\text{dom}})$ iff $g|_{\Sigma_1}, \nu_g \models \varphi_{\text{dom}}(X_1, \dots, X_k)$.

For each $g \in GL(\varphi_{\text{dom}})$ we define the graph $\hat{\tau}(g)$ similar to $\tau(g)$ in Definition 10. The nodes of $\hat{\tau}(g)$ are defined using $g|_{\Sigma_1} \models \varphi_\sigma^c(u, U_1, \dots, U_k)$, where $U_i = \nu_g(X_i)$, rather than $g \models \varphi_\sigma^c(u)$, and similarly for the edges and node labelling of $\hat{\tau}(g)$. The transduction τ is then defined as follows: $\tau = \{ (g|_{\Sigma_1}, \hat{\tau}(g)) \mid g \in GL(\varphi_{\text{dom}}) \}$.

25 Example. Let $m \subseteq \{a\}^* \times \{a, b, \#\}^*$ be the relation

$$\{ (a^n, w\#w) \mid n \geq 0, w \in \{a, b\}^*, |w| = n \}.$$

The relation $\text{ed-gr}(m)$ can be realized by a nondeterministic mso definable transduction, with parameters X_a and X_b . The nodes of the input graph are copied twice, and the parameters determine whether the outgoing edge of a node in the input is copied as a -edge or b -edge, respectively.

The components of the transduction are as follows. The copy set equals $C = \{1, 2\}$, the domain formula $\varphi_{\text{dom}}(X_a, X_b)$ expresses that the input graph is a string representation, and additionally that the sets X_a and X_b form a partition of its nodes.

All input nodes are copied twice: $\varphi_*^1(x, X_a, X_b) = \varphi_*^2(x, X_a, X_b) = \text{true}$.

The edge labels are changed according to the sets X_a and X_b , additionally the last node of the first copy is connected to the first node of the second copy by an $\#$ -edge:

$$\begin{aligned} \varphi_\sigma^{1,1}(x, y, X_a, X_b) &= \varphi_\sigma^{2,2}(x, y, X_a, X_b) = \text{edge}_a(x, y) \wedge x \in X_\sigma, \\ &\quad \text{for } \sigma = a, b, \\ \varphi_\#^{1,2}(x, y, X_a, X_b) &= \neg(\exists z)\text{edge}_a(x, z) \wedge \neg(\exists z)\text{edge}_a(z, y), \\ \varphi_\sigma^{i,j}(x, y, X_a, X_b) &= \text{false, for all other combinations } i, j, \sigma. \end{aligned}$$

Mapping aaa to $abb\#abb$ can be realized by taking the valuation $\nu(X_a) = \{1\}$, $\nu(X_b) = \{2, 3, 4\}$.

Note that this example can be changed such that it uses only one parameter, as the sets represented by the parameters are complementary. \square

We use grNMSO , NMSOS_{nd} , and NMSOS to denote the nondeterministic counterparts of the families grMSO , MSOS_{nd} , and MSOS , respectively. The family of (nondeterministic) 2gsm transductions is denoted by 2NGSM .

Unlike the deterministic case, the power of the nondeterministic 2gsm is incomparable to that of the nondeterministic mso definable string transduction. First, because the number of parameter valuations is finite, every nondeterministic mso transduction is finitary. This is not true for the 2gsm, which can realize the (non-finitary) transduction $\{(a^n, a^{mn}) \mid m, n \geq 1\}$, by nondeterministically choosing the number m of copies made of the input.

On the other hand, the nondeterministic mso transduction of the previous example cannot be realized by a 2gsm.

26 Lemma. *Let $m \subseteq \{a\}^* \times \{a, b, \#\}^*$ be the relation $\{(a^n, w\#w) \mid n \geq 0, w \in \{a, b\}^*, |w| = n\}$. Then $m \notin 2\text{NGSM}$.*

Proof. Assume m is realized by a (nondeterministic) 2gsm \mathcal{M} with k states. Choose n such that $2^n > k \cdot (n + 2)$. Consider the behaviour of \mathcal{M} on input a^n . The input tape, containing $\vdash a^n \dashv$, has $n + 2$ positions. Hence, \mathcal{M} has $k \cdot (n + 2)$ configurations on this input. Consider the configuration assumed by \mathcal{M} when it has just written the symbol $\#$ on its output tape. As there are 2^n possible output strings $w\#w$ for a^n , there exist two strings w_1 and w_2 for which this configuration is the same. This means that we can switch the computation of $(a^n, w_1\#w_1)$ halfway to the computation of $(a^n, w_2\#w_2)$ obtaining a computation for $(a^n, w_1\#w_2)$ with $w_1 \neq w_2$, which is not an element of m . \square

It is not difficult to see that the relation m from the lemma, can be realized by the composition of two 2gsm's, the first nondeterministically mapping a^n to a string $w \in \{a, b\}^*$ with $|w| = n$, the second (deterministically) doubling its input w to $w\#w$. This shows that 2NGSM is not closed under composition, as proved in [Kie75] for the corresponding families of output languages. In fact, the families 2NGSM^k of compositions of k 2gsm transductions form a strict hierarchy, as proved in [Gre78c, Eng82, Eng91b] (again for the corresponding families of output languages).

However, the nondeterministic mso transductions are closed under composition [Cou97, Prop. 5.5.6].

27 Proposition. *grNMSO, and consequently NMSOS and NMSOS_{nd} are closed under composition.*

By grREL we denote the family of (nondeterministic) node relabellings for graphs. A relation in $\text{GR}(\Sigma_1, \Gamma) \times \text{GR}(\Sigma_2, \Gamma)$ is a *node relabelling* if there exists a relation $R \subseteq \Sigma_1 \times \Sigma_2$ such that the images of a graph g are exactly those graphs that can be obtained from g by replacing every occurrence of a node label σ by an element of $R(\sigma)$, leaving edges and their labels unchanged.

We use REL to denote the family of (nondeterministic) *string relabellings*, related to grREL through the mapping nd-gr.

We observe the following elementary relationship between deterministic and nondeterministic mso definable graph transductions.

28 Theorem. $\text{grNMSO} = \text{grREL} \circ \text{grMSO}$.

Proof. The proof of the first inclusion $\text{grNMSO} \subseteq \text{grREL} \circ \text{grMSO}$ is implicit in our definition of grNMSO . The nondeterminism of an mso transduction τ with parameters X_1, \dots, X_k can be ‘pre-processed’ by a relabelling ρ that maps each node label $\sigma \in \Sigma_1$ nondeterministically to a symbol $(\sigma, f) \in \Sigma_1 \times \{0, 1\}^\Xi$, where $\Xi = \{X_1, \dots, X_k\}$. The valuation of X_i has now become a part of the labelling, and we change the domain formula $\varphi_{\text{dom}}(X_1, \dots, X_k)$, the node formulas $\varphi_\sigma^c(x, X_1, \dots, X_k)$, and the edge formulas $\varphi_\gamma^{c_1, c_2}(x, y, X_1, \dots, X_k)$ that specify the mso transduction accordingly. Each atomic subformula $y \in X_i$ in such a formula is replaced by the disjunction $\bigvee_{f(X_i)=1, \sigma \in \Sigma_1} \text{lab}_{(\sigma, f)}(y)$, and each atomic subformula $\text{lab}_\sigma(y)$ is replaced by $\bigvee_{f: \Xi \rightarrow \{0, 1\}} \text{lab}_{(\sigma, f)}(y)$. In this way we obtain ‘deterministic’ equivalents $\hat{\varphi}_{\text{dom}}, \hat{\varphi}_\sigma^c(x), \hat{\varphi}_\gamma^{c_1, c_2}(x, y)$ for mso transduction $\hat{\tau}$. We now have $\tau = \rho \circ \hat{\tau}$ which follows by observing that for a graph $g \in \text{GR}(\Sigma_1 \times \{0, 1\}^\Xi, *)$, $g \models \hat{\varphi}_{\text{dom}}$ if and only if $g|_{\Sigma_1}, \nu_g \models \varphi_{\text{dom}}(X_1, \dots, X_k)$, and similarly for the other formulas.

For the converse inclusion $\text{grNMSO} \supseteq \text{grREL} \circ \text{grMSO}$, it suffices to note that each nondeterministic node relabelling is a nondeterministic mso definable graph transduction. The inclusion then follows from the closure of grNMSO under composition, Proposition 27.

Let $R \subseteq \Sigma_1 \times \Sigma_2$ define a graph node relabelling. We formalize it as mso graph transduction from $\text{GR}(\Sigma_1, \Gamma)$ to $\text{GR}(\Sigma_2, \Gamma)$ by choosing parameters $X_\tau, \tau \in \Sigma_2$, with the intended meaning that a node belonging to X_τ will be relabelled into τ .

The domain formula φ_{dom} expresses that the X_τ form an ‘admissible’ parameter set by demanding each node to be in exactly one of the X_τ , and additionally, if a node has label σ , then X_τ containing this node satisfies $\tau \in R(\sigma)$:

$$(\forall x) \bigvee_{\tau \in \Sigma_2} (x \in X_\tau \wedge \bigwedge_{\tau' \neq \tau} x \notin X_{\tau'}) \wedge (\forall x) \bigwedge_{\sigma \in \Sigma_1} (\text{lab}_\sigma(x) \rightarrow \bigvee_{\tau \in R(\sigma)} x \in X_\tau)$$

Each node is copied once, relabelled according to X_τ :

$$\begin{aligned} C &= \{1\}, \\ \varphi_\tau^1 &= x \in X_\tau, \tau \in \Sigma_2, \\ \varphi_\gamma^{1,1} &= \text{edge}_\gamma(x, y), \gamma \in \Gamma. \end{aligned}$$

□

As we have observed, any string relabelling can be ‘lifted’ to a graph node relabelling using the graph interpretation nd-gr of strings. By restricting the previous result to those graph transductions that result from strings, we obtain a result for mso definable string transductions in the node interpretation.

29 Corollary. $\text{NMSOS}_{\text{nd}} = \text{REL} \circ \text{MSOS}_{\text{nd}}$.

In addition to REL , we need MREL denoting the family of *marked string relabellings*, that map a string w first to the ‘marked version’ $\vdash w \dashv$, and then apply a string relabelling.

30 Theorem. $\text{NMSOS} = \text{MREL} \circ \text{MSOS}$.

Proof. First, the inclusion from left to right. Let $m \in \text{NMSOS}$, i.e., $\text{ed-gr}(m) \in \text{grNMSO}$.

Consider the string transduction $m' = \{ (\vdash w \dashv, z) \mid (w, z) \in m \}$. Then m' is an element of NMSOS_{nd} , as $\text{nd-gr}(m')$ equals the composition $\text{tape}(\text{id}) \circ \text{ed-gr}(m) \circ \text{ed2nd}$ of (nondeterministic) mso definable graph transductions, where $\text{tape}(\text{id})$ is the mapping from $\text{nd-gr}(\vdash w \dashv)$ to $\text{ed-gr}(w)$, cf. Lemma 22.

By the corollary above, and Lemma 18, $m' \in \text{REL} \circ \text{MSOS}_{\text{nd}} \subseteq \text{REL} \circ \text{MSOS}$. Consequently, as m equals the ‘marking’ from w to $\vdash w \dashv$ followed by m' , $m \in \text{MREL} \circ \text{MSOS}$.

For the reverse inclusion, $\text{NMSOS} \supseteq \text{MREL} \circ \text{MSOS}$, note that every marked relabelling can be decomposed into a marking and a relabelling, each of which we will show to be a (nondeterministic) mso transduction. The inclusion then follows from the closure of NMSOS under composition.

The marking mapping w to $\vdash w \dashv$ is easily seen to be an element of MSOS , either by direct construction, or by constructing a 2dgsms for that task, and applying Theorem 23.

Finally, to show that $\text{REL} \subseteq \text{NMSOS}$ one closely follows the argumentation in the proof of $\text{grREL} \subseteq \text{grNMSO}$, Theorem 28. As we relabel edges, rather than nodes, in the representation $\text{ed-gr}(w)$ of a string w , but still have parameters ranging over nodes, we use the parameters for the source node of an edge to determine the new label of its outgoing edge (cf. Example 25): φ_{dom} is as before, but we now have $\varphi_*^1 = \text{true}$, and $\varphi_\tau^{1,1} = \text{edge}(x, y) \wedge (x \in X_\tau)$. \square

For completeness we note that the above result cannot be strengthened to $\text{NMSOS} = \text{REL} \circ \text{MSOS}$, as the relations on the right side are functional for the empty string λ . This is not necessarily true for NMSOS.

31 Example. The string transduction $\{(\lambda, a), (\lambda, b)\}$ in $a^* \times \{a, b\}^*$ is realized by the following nondeterministic mso transduction, in the edge representation. The single parameter X determines whether λ is mapped to a or to b . Let

$$\begin{aligned} \varphi_{\text{dom}} &= (\exists x)(\forall y)(y = x) \wedge \neg \text{edge}_a(x, x), \\ C &= \{1, 2\}, \\ \varphi_*^1 &= \varphi_*^2 = \text{true}, \\ \varphi_\sigma^{1,1} &= \varphi_\sigma^{2,1} = \varphi_\sigma^{2,2} = \text{false}, \text{ for } \sigma \in \{a, b\}, \text{ and} \\ \varphi_a^{1,2} &= x \in X, \\ \varphi_b^{1,2} &= \neg(x \in X). \end{aligned}$$

□

Combining the previous two results (that relate the nondeterministic and deterministic mso transductions) with the equalities between deterministic mso transductions and deterministic gsm mappings of Theorem 23, we directly obtain the following result.

32 Theorem.

$$\text{NMSOS} = \text{MREL} \circ 2\text{DGSM} \quad \text{and} \quad \text{NMSOS}_{\text{nd}} = \text{REL} \circ 2\text{DGSM}\lambda.$$

6 Finite Visit Machines

Rajlich [Raj75] observes that 2gsm are more powerful than 2dgs (as generative devices, by considering their output languages, i.e., the ranges of the transductions). He demonstrates that this is mainly due to the ability of the 2gsm to visit each of the positions of its input an unbounded number of times.

Motivated by this result, we consider transducers that have a fixed bound on the number of times they visit each of their input positions –we call this the finite visit property– and relate these to the (nondeterministic) mso transductions.

We show that the nondeterministic mso definable string transductions are exactly those transductions that are realized by the composition of two 2gsm with the finite visit property. Note that one direction of this result follows from Theorem 32.

Moreover, we characterize the nondeterministic mso definable string transductions as those compositions of 2gsm’s that realize finitary transductions, i.e., transductions that define a finite number of images for every input string.

A more direct characterization can be obtained by considering 2gsm that are allowed to rewrite the symbols on their input tape (but with the finite visit property). These machines exactly match the mso definable string transductions, both in the deterministic case and the nondeterministic case.

The finite visit property was studied in, e.g., [Hen65, Raj75, Gre78a, Gre78b, Gre78c, ERS80, Eng82].

6.1 Finite visit two-way generalized sequential machines

A computation of a 2gsm is called *k-visiting* if each of the positions of the input tape is visited at most k times. The 2gsm \mathcal{M} is called *finite visit*, if there is a constant k such that, for each pair (w, z) in the transduction realized by \mathcal{M} , there exists a k -visiting computation for (w, z) . The family of string transductions realized by finite visit nondeterministic 2gsm is denoted by $2\text{NGSM}_{\text{fin}}$.

Note that our definition is rather weak, as the machine may have many computations that are not k -visiting, either without any chance of reaching the final state, or with loops in the computation that produce no output.

If a deterministic 2gsm visits a position of the input tape twice in the same state, then the computation will enter an infinite loop that will not reach the

final state. This implies the well-known fact that every deterministic 2gsm is finite visit, where we choose for k the number of states of the machine. A similar argument enables us to prove the following characterization of finite visit transductions in terms of transductions that map each input string into a finite number of output strings.

33 Lemma. *Let m be a string transduction. Then $m \in 2\text{NGSM}_{\text{fin}}$ iff $m \in 2\text{NGSM}$ and m is finitary.*

Proof. Clearly, the length of the output of a k -visiting computation on input w is at most k times the length of $\vdash w \dashv$. Hence the implication from left to right.

As for the other implication, assume that the finitary transduction m is realized by a 2gsm \mathcal{M} . If during a (successful) computation for $(w, z) \in m$, \mathcal{M} visits the same position twice in the same state, then it did not write symbols to the output in the meantime, because otherwise \mathcal{M} has infinitely many output strings for the present input, as an easy pumping argument shows. Hence we may omit this excursion from the computation. Consequently, there is a computation of \mathcal{M} for (w, z) that does not visit each of the tape positions more than k times, where k is the number of states of \mathcal{M} . Hence \mathcal{M} *itself* is finite visit. \square

It is well known (see, e.g., [Fis69, ChJá77, Gre78a, Gre78b, AhU170]) that the computation of a finite visit 2gsm on an input tape can be coded as a string of ‘visiting sequences’ (strongly related to ‘crossing sequences’, cf. [Rab63, Hen65, HoU179, Bir96]). We recall how this can be done, without going into details.

We consider several types of visits during a computation, differing in the direction (-1 , 0 , or $+1$) of the steps taken by the machine just before and just after the visit. Additionally, a visit may be either the first or the last visit of the computation.

Given a computation of a 2gsm, the *visiting sequence* of a position of the input tape is the sequence that starts with the symbol σ on the tape, followed by the consecutive visits of the machine to that position. Each of the visits is given as a 4-tuple $(\neg\epsilon, p, +\epsilon, \alpha)$ consisting of the direction $\neg\epsilon$ of the move before the visit, the state p during the visit, the direction $+\epsilon$ of the move after the visit, and the string α written to the output during that move. For the first visit we take $\neg\epsilon = *$, for the last visit we take $+\epsilon = *$.

We illustrate this notion with an example.

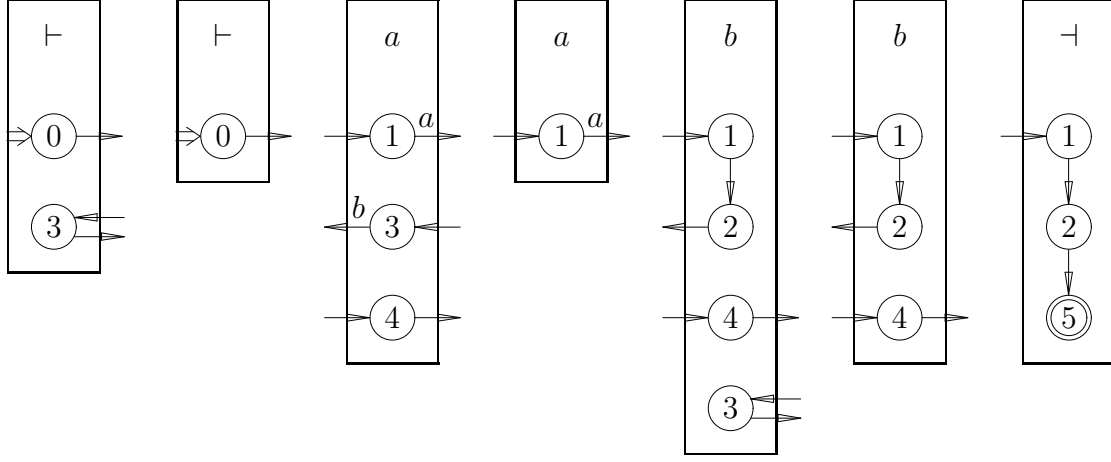


Figure 6: Visiting sequences for Example 2, cf. Example 34.

34 Example. Consider the 2dgs from Example 2. Each of the visiting sequences during a successful computation is one of the following.

$$\begin{aligned}
& \langle \vdash, (*, 0, +1, \lambda), (-1, 3, +1, \lambda) \rangle \\
& \langle \vdash, (*, 0, +1, \lambda) \rangle \\
& \langle a, (+1, 1, +1, a), (-1, 3, -1, b), (+1, 4, +1, \lambda) \rangle \\
& \langle a, (+1, 1, +1, a) \rangle \\
& \langle b, (+1, 1, 0, \lambda), (0, 2, -1, \lambda), (+1, 4, +1, \lambda), (-1, 3, +1, \lambda) \rangle \\
& \langle b, (+1, 1, 0, \lambda), (0, 2, -1, \lambda), (+1, 4, +1, \lambda) \rangle \\
& \langle \neg, (+1, 1, 0, \lambda), (0, 2, 0, \lambda), (0, 5, *, \lambda) \rangle
\end{aligned}$$

These visiting sequences are depicted in a suitable graphical manner in Figure 6, cf. Figure 1. \square

Each visiting sequence must satisfy some syntactical constraints.

First, the directions of the visits are ‘alternating’. This means that the first visit enters from the left ($\neg\epsilon = +1$, with the exception for $\sigma = \vdash$ which starts in the initial state with $\neg\epsilon = *$); then, if the move after the i -th visit equals $\neg\epsilon = -1, 0, +1$, then the move prior to the $i+1$ -st visit to the same position must equal $\neg\epsilon' = +1, 0, -1$, respectively. Only the last visit of a

sequence can have ${}^+\epsilon = *$, in case the state is final, signalling the end of a computation.

Secondly, the direction ${}^+\epsilon$ of the move after the visit, and the string α written to the output, must correspond to an instruction of the machine for the given input symbol σ and the given state p . Additionally, when ${}^+\epsilon = 0$, the new state given by the instruction must match the next visit of the visiting sequence.

Clearly, also neighbouring visiting sequences for a given computation must satisfy several constraints. If a visiting sequence has k ‘crossings’ to the right, either outgoing visits $(\neg\epsilon, p, +1, \alpha)$ or incoming visits $(-1, p, {}^+\epsilon, \alpha)$ —they alternate—then the visiting sequence to the right has exactly k matching crossings to the left, matching both in direction (which implicitly follows from the restrictions on single visiting sequences above) and in state change for the machine. Note that a visit $(-1, p, +1, \alpha)$ represents two crossings.

Finally, the first visiting sequence of a computation should start with a visit $(*, q_{in}, {}^+\epsilon, \alpha)$, and exactly one visiting sequence should end with a visit $(\neg\epsilon, q_f, *, \alpha)$.

When we bound the number of visits to each position, the visiting sequences come from a finite set, and we can interpret these sequences as symbols from a finite alphabet. Each k -visiting computation is specified by a string over this alphabet, and we will call these strings *k-tracks*. (E.g., the track in Figure 7 specifies the computation of the 2dgsm of Example 2 on input a^3b^2aba , cf. Figure 1). It should be obvious from the above remarks that the language of such specifications is regular (see, e.g., Lemma 2.2 of [Gre78a], or Lemma 1 of [ChJá77]). For instance, it is the heart of the proof in [HoU179, Theorem 2.5] of the result that two-way finite state automata are equivalent to their one-way counterparts [RaSc59, She59].

35 Proposition. *Let \mathcal{M} be a 2gsm, and let k be a constant. The k -tracks for successful k -visiting computations of \mathcal{M} form a regular language.*

From this result, using standard techniques (see e.g., [ChJá77, Lemma 1]) we obtain the following decomposition of finite visit nondeterministic 2gsm transductions. Note that this decomposition already features in Theorem 32 as characterization of NMSOS.

36 Lemma. $2\text{NGSM}_{\text{fin}} \subseteq \text{MREL} \circ 2\text{DGSM} = \text{NMSOS}.$

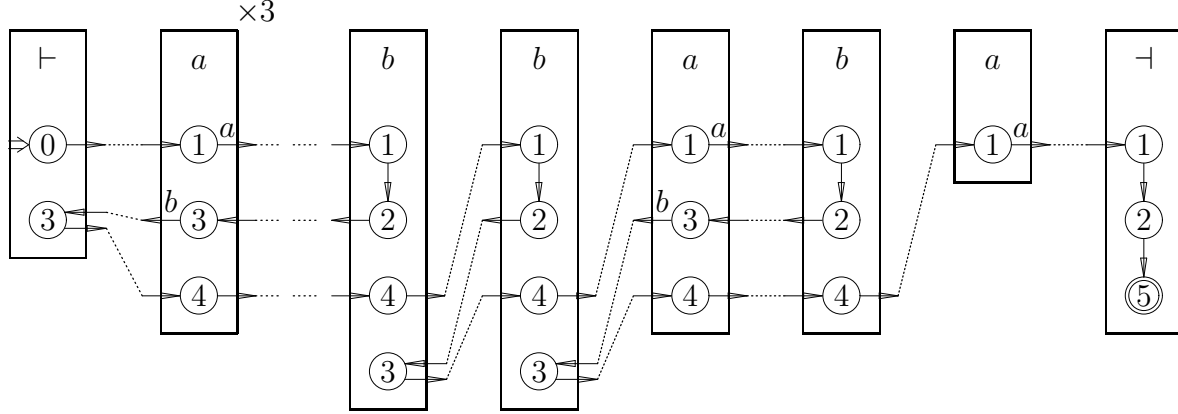


Figure 7: Track for $\vdash a^3 b^2 a b a \dashv$, Example 2.

Proof. Let \mathcal{M} be a 2gsm, finite visit for constant k ; each pair (w, z) in the transduction realized by \mathcal{M} can be computed by a k -visiting computation.

We may decompose the behaviour of \mathcal{M} on input w as follows. First, a relabelling of $\vdash w \dashv$ guesses a string of k -visiting sequences, one for each position of the input tape. Then, a 2dgsm verifies in a left to right scan whether the string specifies a valid computation, a track, of \mathcal{M} for w , cf. Proposition 35. If this is the case, the 2dgsm returns to the left tape marker \vdash and simulates \mathcal{M} on this input, following the k -visiting computation previously guessed.

When changing from one tape position to a neighbouring position, the 2dgsm records the ‘crossing number’ of that move, i.e., the number of times it crossed the border between these two tape positions (in one direction or another). The crossing number can be read by inspecting the directions of the moves stored in the visiting sequence. It is used to ‘enter’ the next visiting sequence at the right visit, cf. Figure 7. \square

37 Theorem. $\text{NMSOS} = 2\text{NGSM}_{\text{fin}} \circ 2\text{NGSM}_{\text{fin}}$.

Proof. By the last lemma, $2\text{NGSM}_{\text{fin}} \subseteq \text{NMSOS}$. As the right-hand side of this inclusion is closed under composition (Proposition 27) we have the inclusion $2\text{NGSM}_{\text{fin}} \circ 2\text{NGSM}_{\text{fin}} \subseteq \text{NMSOS}$.

According to Theorem 32, NMSOS equals $\text{MREL} \circ 2\text{DGSM}$. The inclusion from left to right follows from the fact that both $\text{MREL} \subseteq 2\text{NGSM}_{\text{fin}}$ and $2\text{DGSM} \subseteq 2\text{NGSM}_{\text{fin}}$. \square

It is instructive to note that this characterization implies the (apparently new) result that $2\text{NGSM}_{\text{fin}} \circ 2\text{NGSM}_{\text{fin}}$ is closed under composition. This should be contrasted to the fact that $2\text{NGSM}_{\text{fin}}$ itself is not closed under composition. This follows from the observation from the preceding section, that the relation m from Example 25 does not belong to $2\text{NGSM} \supseteq 2\text{NGSM}_{\text{fin}}$ (Lemma 26). As we have observed, it can be realized as combination of two 2gsm 's, the first one nondeterministically changing a string a^n to a string $w \in \{a, b\}^*$ with $|w| = n$, the second one duplicating w into $w\#w$. Both of these 2gsm 's are finite visit. (Alternatively, by Example 25, $m \in \text{NMSOS}$ which equals $2\text{NGSM}_{\text{fin}}^2$ as we just have seen.)

The families 2DGSM , $2\text{NGSM}_{\text{fin}}$, and $2\text{NGSM}_{\text{fin}}^2$ form a hierarchy of transductions. However, as far as their output languages are concerned (ranges, or equivalently, with regular input) these three families are equally powerful [Kie75, Gre78b].

Recall that the families 2NGSM and NMSOS are incomparable, see the discussion preceding Lemma 26. We have a surprising characterization for their intersection.

38 Theorem. $2\text{NGSM} \cap \text{NMSOS} = 2\text{NGSM}_{\text{fin}}$.

Proof. Obviously $2\text{NGSM}_{\text{fin}} \subseteq 2\text{NGSM}$, while $2\text{NGSM}_{\text{fin}} \subseteq \text{NMSOS}$ by Theorem 37, which proves the inclusion from right to left.

The reverse implication is immediate from Lemma 33: recall that transductions in NMSOS are finitary because the number of parameter valuations is finite. \square

Combining this theorem and the related Lemma 33, we obtain that a 2gsm string transduction is mso definable if and only if it is finitary. This generalizes a similar result of Courcelle [Cou94, Proposition 6.1] for rational transductions (i.e., string transductions realized by 2gsm never moving to the left). It can be extended to arbitrary compositions of two-way gsm 's, as we shall see in our next main result, Theorem 42.

As a preparation to this result (and its proof) we like to point out that ‘pumping’ computations for finite visit transductions (iterating suitable segments of tracks) does not only result in duplication of parts of the output, but may also rearrange neighbouring segments of the output. We illustrate this with an example.

39 Example. The 2gsm \mathcal{M} has states 1 to 6, initial state 1, final state 6, and transitions $(p, \sigma, q, \alpha, \epsilon, p, \lambda, 0)$ where the move q, α, ϵ for each pair $p \in \{1, 2, \dots, 5\}$, $\sigma \in \{\vdash, a, b, \dashv\}$ is given in the following matrix.

	1	2	3	4	5
\vdash	$1, \lambda, +1$	$3, b, 0$	$3, \lambda, +1$	$5, b, 0$	$5, \lambda, +1$
a	$1, a, +1$	$2, a, -1$	$3, a, +1$	$4, a, -1$	$5, a, +1$
b	$2, \lambda, -1$	$4, \lambda, -1$	$1, \lambda, +1$	$5, \lambda, +1$	$3, \lambda, +1$
\dashv	$2, c, 0$	$2, \lambda, -1$	$4, c, 0$	$4, \lambda, -1$	$6, \lambda, 0$

(Note that the machine is nondeterministic in our setting, but is obtained by adding dummy alternatives to a deterministic automaton in the 5-tuple framework, see Section 2.)

On each segment of a 's of the input \mathcal{M} makes five passes in states 1 to 5, each in alternate directions, while copying the letters to the output.

On a letter b the machine does not generate output, but it performs a permutation of the order in which the two neighbouring segments of a 's are read. This is best explained by looking at the computations on the input strings $a^3b^i a^2$, $i = 0, 1, 2$ as depicted in Figure 8. The output strings for these inputs are given in the following table.

input string	output string
$a^5 = a^3b^0a^2$	$a^5ca^5ba^5ca^5ba^5 = a^3(a^2ca^2)(a^3ba^3)(a^2ca^2)(a^3ba^3)a^2$
$a^3b^1a^2$	$a^6ba^5ca^5ba^5ca^4 = a^3(a^3ba^3)(a^2ca^2)(a^3ba^3)(a^2ca^2)a^2$
$a^3b^i a^2, i \geq 2$	$a^6ba^6ba^5ca^4ca^4 = a^3(a^3ba^3)(a^3ba^3)(a^2ca^2)(a^2ca^2)a^2$

As we have seen, the introduction of the symbol b in the input does not generate new output. Instead, it rearranges the parts of the computation that extend to both sides of the symbol.

Consider the boundary between two tape positions, where we want to insert a symbol b . Let $x_1, z_1, x_2, z_2, x_3, z_3$ be the strings written to the output during the consecutive parts of the computation that visit the left (x_i) and right (z_i) segments of the tape, see Figure 9. The output generated is thus $x_1z_1x_2z_2x_3z_3$.

Now, we introduce b at the selected boundary, and obtain the new output $x_1x_2z_1x_3z_2z_3$. This rearrangement of the output can be formalized by the application of the substitution $\sigma_b : [z_1, z_2, z_3 \leftarrow \lambda, z_1, z_2z_3]$ – where z_i is a formal parameter rather than a specific string.

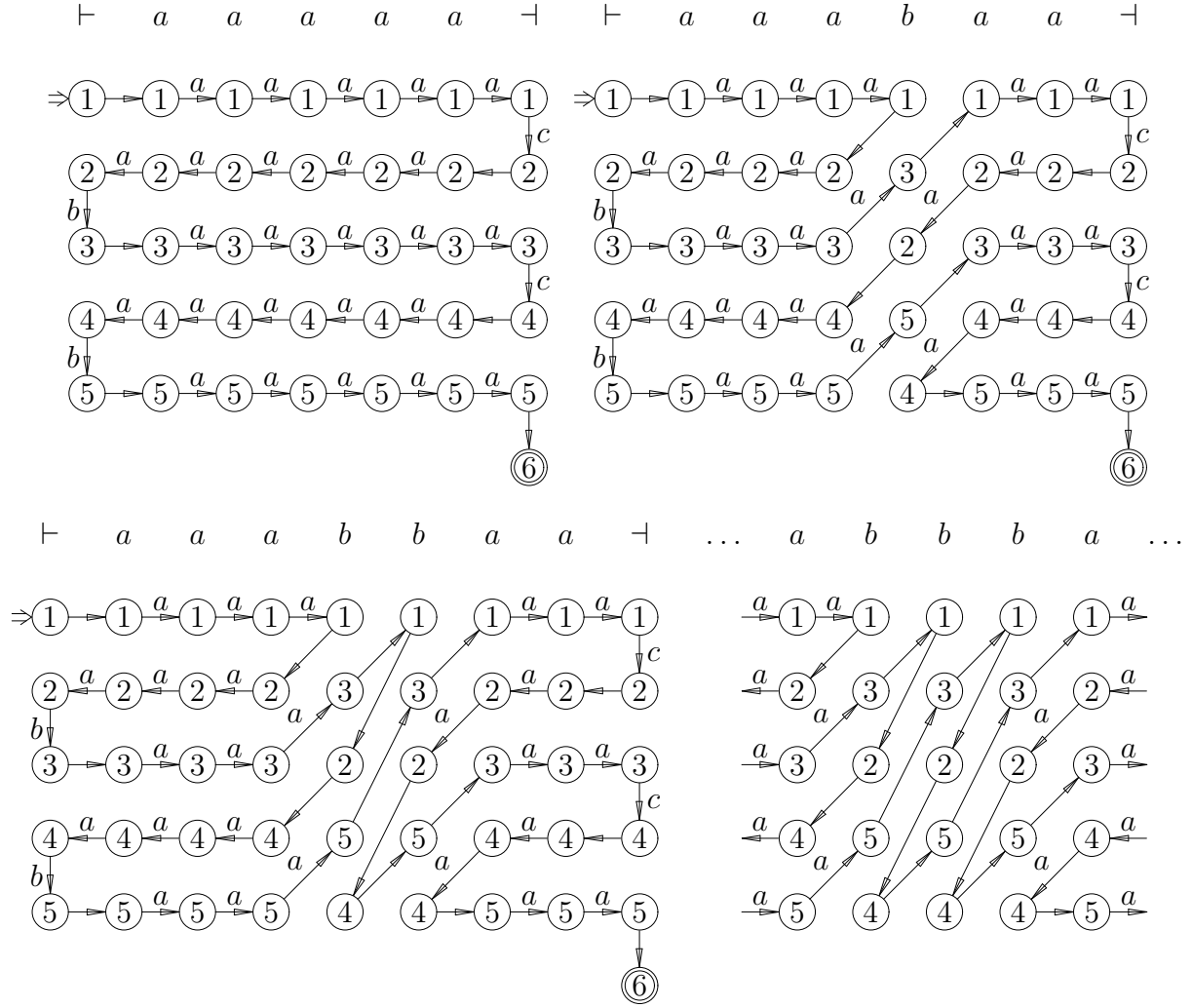


Figure 8: Computations for Example 39

The effect of introducing bb can be computed by the composition $\sigma_b \sigma_b : [z_1, z_2, z_3 \leftarrow \lambda, \lambda, z_1 z_2 z_3]$, which defines the rearrangement $x_1 x_2 x_3 z_1 z_2 z_3$ of the output. Note that $\sigma_b^i = \sigma_b^2$ for $i \geq 2$. \square

40 Lemma. *Let m be a finitary string transduction, and let X be a family of string transductions.*

If $m \in X \circ \text{2NGSM} \circ \text{2DGSM}$, then $m \in X \circ \text{MREL} \circ \text{2DGSM}$.

Proof. Assume that the finitary transduction m is a composition $m = m_0 \circ m_1 \circ m_2$ as in the statement of the lemma; $m_0 \in X$, m_1 realized by the 2gsm \mathcal{M}_1 , and m_2 realized by the 2dgsm \mathcal{M}_2 . As to be expected, the unknown family X will not feature in our arguments, but later will enable us to apply the result in a context. In fact, we show how to replace $m_1 \circ m_2$ by $\check{m}_1 \circ \check{m}_2 \in \text{MREL} \circ \text{2DGSM}$ such that $m_0 \circ m_1 \circ m_2 = m_0 \circ \check{m}_1 \circ \check{m}_2$. Hence $m_1 \circ m_2$ equals $\check{m}_1 \circ \check{m}_2$ on the range of m_0 .

Reconsider the proof of Lemma 36, where a k -visit 2gsm is decomposed into a relabelling that guesses a k -visiting sequence for each position of the input tape, and a 2dgsm that verifies in a single left-to-right pass whether the resulting string defines a k -track, and then deterministically simulates the specified computation for the original input. Alternatively, by combining the verification phase with the relabelling, we may decompose the k -visit 2gsm into a one-way gsm that nondeterministically writes a k -track, and a 2dgsm simulating the computation.

We apply that new decomposition to \mathcal{M}_2 , and immediately observe that the first phase (guessing and writing a track) can be performed by \mathcal{M}_1 using a straightforward direct product construction.

Summarizing: we have replaced the composition $m_1 \circ m_2$ by a new composition $m'_1 \circ m'_2$ realized by \mathcal{M}'_1 followed by \mathcal{M}'_2 , where \mathcal{M}'_1 is a 2gsm that writes valid tracks for the 2dgsm \mathcal{M}'_2 . Let \mathcal{M}'_2 be k -visit.

We continue by demonstrating that we need not consider all computations of \mathcal{M}'_1 , instead it suffices to put a bound on the number of visits that the machine makes to each of the positions of its input. This will change the transduction m'_1 realized by \mathcal{M}'_1 , but not the composition $m_0 \circ m'_1 \circ m'_2$ (due to m being finitary).

Consider the behaviour of \mathcal{M}'_1 on input w , where w is in the range of m_0 . Fix a position on the tape $\vdash w \dashv$ and a state of \mathcal{M}'_1 , and split the output of \mathcal{M}'_1 during the computation into segments, corresponding to the consecutive

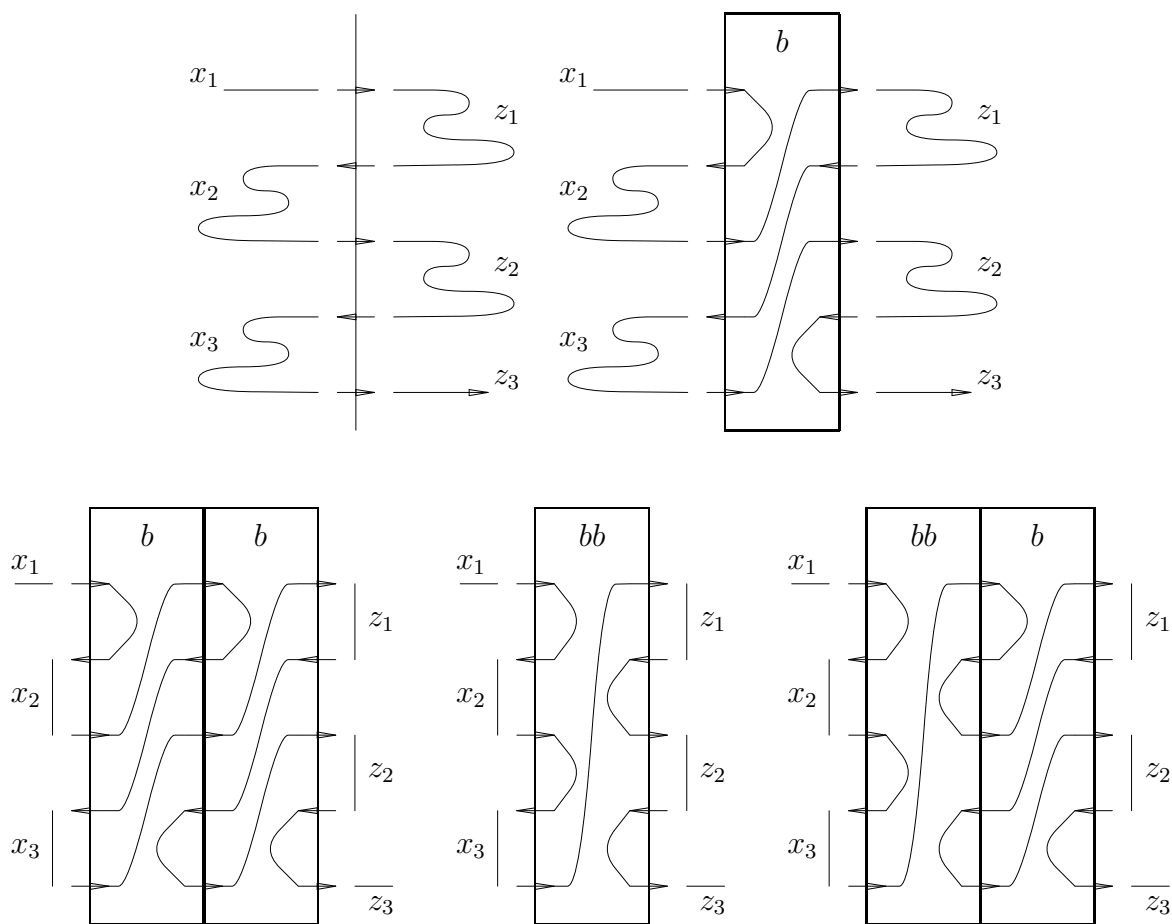


Figure 9: Visualization of rearrangements

visits to the selected position in the selected state. \mathcal{M}'_1 writes $xy_1y_2\cdots y_tz$ where y_i is written during the excursions in between consecutive visits. We assume $t \geq 1$.

Returning to the same position and state, each of the excursions can be repeated in (or omitted from) the computation of \mathcal{M}'_1 , so the machine may produce every string xyz , $y \in \{y_1, y_2, \dots, y_t\}^*$ as possible output on input w . By our previous construction, each output of \mathcal{M}'_1 forms a k -track for the second machine \mathcal{M}'_2 . This implies that \mathcal{M}'_2 does not generate output during any of its visits to the segments y_i , as m is supposed to be finitary.

At first glance, the excursion of \mathcal{M}'_1 writing $y = y_1 \cdots y_t$ can be omitted: the second machine \mathcal{M}'_2 does not generate output when it visits the segment y during its simulation of the specified computation. However, the previous example shows that y (or in fact any segment y_i) may have its effect on the output of \mathcal{M}'_2 by rearranging parts of the adjacent computation that leave the segment y (to the left or to the right) in order to return there later.

We consider the computation of \mathcal{M}'_2 specified by the track xyz from the viewpoint of the segment y . Starting from the leftmost position of x , the computation enters y from the left. Before leaving the segment for the last time, the computation makes several tours outside y .

Such a tour of \mathcal{M}'_2 to the left of the segment y , in x , corresponds to two consecutive visits $(-\epsilon, p, -1, \lambda)$ and $(+1, p', +\epsilon', \lambda)$ in the first visiting sequence of y , meaning the computation leaves the segment to the left in state p , returning there later in state p' . A symmetric observation holds for tours to the right, in z , and consecutive visits in the last visiting sequence of y .

Hence, the relative order of those tours that leave to the left is fixed by the last visiting sequence of x , similarly for the tours to the right. The relative order of all tours (left and right taken together) is determined by the segment y . Replacing y by another string in $\{y_1, y_2, \dots, y_t\}^*$ will not change the tours in x and z , but it may rearrange the relative order of tours to the left and tours to the right.

A visiting sequence for \mathcal{M}'_2 contains at most k visits. Hence, there are less than k tours to each side of the segment. Together these at most $2k$ tours may be ordered in less than $\kappa = \binom{2k}{k}$ ways (the orders of the tours at the same side of the segment are fixed).

Now we are able to apply a pumping argument to the segment $y = y_1 \cdots y_t$. If $t > \kappa$, then two of the prefixes $y_1 \cdots y_{i_1}$, $y_1 \cdots y_{i_2}$, $i_1 < i_2$, define the same rearrangement on the adjacent tours, and thus we may re-

place $y_1 \cdots y_{i_2}$ by $y_1 \cdots y_{i_1}$ in the output xyz of \mathcal{M}'_1 . The resulting track $xy_1 \cdots y_{i_1} y_{i_2+1} \cdots y_t z$ defines a computation for \mathcal{M}'_2 that results in the same output as the original track xyz . Thus, we may assume that $t \leq \kappa$.

Consequently, we allow for all possible rearrangements, and hence for all possible outputs of \mathcal{M}'_2 , by taking κ as the bound on the number of visits of \mathcal{M}'_1 to a fixed position in a fixed state.

Now that we have limited the number of visits of \mathcal{M}'_1 to κ times the size of its state set, we can replace \mathcal{M}'_1 by a decomposition in $\text{MREL} \circ 2\text{DGSM}$, using again the argumentation of Lemma 36. Thus, $m'_1 \circ m'_2$ is replaced by a composition in $(\text{MREL} \circ 2\text{DGSM}) \circ 2\text{DGSM}$. The result follows, as 2DGSM is closed under composition, Proposition 4. \square

The variable family X in the previous result allows us to apply the lemma in the context of an arbitrary sequence of 2gsm transductions.

41 Theorem. *Let m be a string transduction, and let $k \geq 1$. If $m \in 2\text{NGSM}^k$, and m is finitary, then $m \in \text{MREL} \circ 2\text{DGSM}$.*

Proof. Observe that $2\text{NGSM} \circ \text{MREL} \subseteq 2\text{NGSM}$ by an obvious construction.

Let $k \geq 1$. Assume that $m \in 2\text{NGSM}^k \circ 2\text{DGSM}$ is finitary. We have by the previous lemma, $m \in 2\text{NGSM}^{k-1} \circ \text{MREL} \circ 2\text{DGSM}$, which equals $2\text{NGSM}^{k-1} \circ 2\text{DGSM}$ for $k > 1$ (and which equals $\text{MREL} \circ 2\text{DGSM}$ for $k = 1$).

Hence, by induction on k , $m \in 2\text{NGSM}^k \circ 2\text{DGSM}$ implies $m \in \text{MREL} \circ 2\text{DGSM}$, for a finitary string transduction m . As $2\text{NGSM}^k \circ 2\text{DGSM} \supseteq 2\text{NGSM}^k$, the theorem follows. \square

42 Theorem. *Let m be a string transduction. Then $m \in \text{NMSOS}$ iff $m \in \bigcup_{k \geq 1} 2\text{NGSM}^k$ and m is finitary.*

Proof. By Theorem 37, $\text{NMSOS} = 2\text{NGSM}_{\text{fin}}^2 \subseteq 2\text{NGSM}^2$. Additionally, elements of NMSOS are necessarily finitary. This proves the implication from left to right. The reverse implication follows from the last result and the characterization $\text{NMSOS} = \text{MREL} \circ 2\text{DGSM}$ from Theorem 32. \square

It is shown in [Eng82, Theorem 4.9] that every functional transduction in $\bigcup_{k \geq 1} 2\text{NGSM}^k$ is in 2DGSM. Together with Theorem 23 (MSOS = 2DGSM) this gives the following counterpart of Theorem 42.

43 Theorem. *Let m be a string transduction. Then $m \in \text{MSOS}$ iff $m \in \bigcup_{k \geq 1} 2\text{NGSM}^k$ and m is functional.*

A Venn diagram is given in Figure 10, page 56. It illustrates the results from Lemma 33, and Theorems 38, 42, and 43.

6.2 Hennie machines

Extending a finite visit 2gsm with the possibility to rewrite the contents of the cell of the input tape that it is visiting, we obtain the *Hennie machine*, introduced in [Hen65] as an accepting device, and considered as transducer in [Raj75] (under the name ‘bounded crossing transducer’). Alternatively, a Hennie machine is a linear bounded automaton (as transducer, so equipped with a one-way output tape) that is finite visit. We find it, somewhat disguised, in [Gre78b] as ‘one way finite visit preset Turing machine’, where the ‘preset working tape’ should be interpreted as input tape, and the ‘one way input tape’ as output tape.

It should be clear how to extend our basic 2sm model to allow for writing on the input tape, thus we will refrain from giving the full 10-tuple formalization. The families of string transductions realized by nondeterministic and deterministic Hennie machines are denoted by NHM and DHM, respectively.

44 Example. Once again consider our running nondeterministic example (cf. Example 25)

$$m = \{ (a^n, w\#w) \mid n \geq 0, w \in \{a, b\}^*, |w| = n \}.$$

It can be realized by a Hennie machine moving in two consecutive left-to-right passes over the input. First it nondeterministically rewrites the input a^n into a string w with $|w| = n$, while writing this string to the output tape, then it writes w again to the output, copying it from the rewritten input tape. Obviously, the machine is 3-visit. \square

45 Theorem. $\text{NMSOS} = \text{NHM}$.

Proof. In view of Theorem 32 it suffices to prove the equality $\text{NHM} = \text{MREL} \circ 2\text{DGSM}$.

The inclusion of NHM in $\text{MREL} \circ 2\text{DGSM}$ can be proved as Lemma 36, which states this inclusion for $2\text{NGSM}_{\text{fin}}$: the relabelling guesses a string of visiting sequences for the computation of the Hennie machine on the input string; the 2dgsm verifies that this string is a track and simulates the computation. Note that a visiting sequence of a Hennie machine should also record the symbol at the position of the input tape at each visit. It is straightforward to adapt the notions of visiting sequence and k -track in this way, such that Proposition 35 still holds (see [Gre78a, Gre78b, Bir96]).

The reverse inclusion is almost immediate. In two phases the Hennie machine may simulate the composition, first writing the image of the marked relabelling on the tape, and then simulating the 2dgsm on this new tape. There is a minor technicality: for a given input w the initial tape contains $\vdash w \dashv$, and the Hennie machine is supposed to overwrite this string with its relabelling and add two new tape markers (for the simulation of the 2dgsm). Instead, it keeps the relabelling of the tape markers in its finite state memory, rather than overwriting them. \square

Restating the above result as $\text{NHM} = \text{MREL} \circ 2\text{DGSM}$, it generalizes the result of Rajlich [Raj75, Theorem 2.1] that the output languages of non-deterministic Hennie machines equal the output languages of two-way deterministic generalized sequential machines, see also [Gre78a, Thm 2.15(2)].

The above demonstration of the inclusion $\text{MREL} \circ 2\text{DGSM} \subseteq \text{NHM}$ can easily be extended to a proof of $\text{NHM} \circ \text{NHM} \subseteq \text{NHM}$. A Hennie machine can simulate the composition of two of its colleagues by writing the visiting sequences of the first machine onto the input tape. The output tape is contained in this string, conveniently folded over the input tape, ready to be used by the second machine.

We have, however, the closure of NHM under composition for free as a consequence of the above characterization and Proposition 27.

46 Corollary. *NHM is closed under composition.*

In [ChJá77] it is noted that the inclusion $\text{DHM} \circ 2\text{DGSM} \subseteq 2\text{DGSM}$ can be proved analogously to their result that $2\text{DGSM} \circ 2\text{DGSM} \subseteq 2\text{DGSM}$ (i.e., 2DGSM is closed under composition, Proposition 4). That of course implies the equality of the families of transductions realized by deterministic

Hennie machines and those realised by deterministic 2gsm. This equality is rephrased as follows.

47 Theorem. $\text{MSOS} = \text{DHM}$.

Proof. In view of Theorem 23 it suffices to prove the equality $\text{DHM} = \text{2DGSM}$. The inclusion $\text{DHM} \supseteq \text{2DGSM}$ is immediate. We demonstrate the reverse inclusion, much along the lines as sketched in [ChJá77], see also [Eng82, Theorem 4.9].

By Theorem 45, $\text{NHM} = \text{MREL} \circ \text{2DGSM}$. Hence, any Hennie transduction m_H can be decomposed into a marked relabelling ρ and a deterministic 2gsm transduction m_2 . We will argue that for a deterministic Hennie transduction this (nondeterministic) marked relabelling can be realized by a deterministic 2gsm, which shows $\text{DHM} \subseteq \text{2DGSM}$ by the closure of 2DGSM under composition.

Let m_H be a deterministic Hennie transduction, and let $m_H = \rho \circ m_2$ be the decomposition as above. Let w be an input string. As m_H is functional, $m_H(w) = m_2(w')$ for *any* marked relabelling $w' \in \rho(w)$ that belongs to the domain of m_2 . As this domain $\text{dom}(m_2)$ is a regular language [RaSc59, She59], a 2dgsm-rla can be constructed that finds and outputs such a marked relabelling by one pass from left to right over the input, using its look-around to check the remainder of the input for a relabelling of the present input symbol that leads to an element of $\text{dom}(m_2)$. This means that the 2dgsm-rla looks ahead to test the suffix of the tape for membership in the language $\rho^{-1}(L(\mathcal{A}_q))$, where \mathcal{A}_q is a (fixed) one-way deterministic finite state automaton accepting $\text{dom}(m_2)$ except that the initial state is changed to q which is the state where \mathcal{A} would be after reading the output generated by the 2dgsm-rla on the prefix, including the relabelling chosen for the present symbol. \square

Finale. In this section we have obtained a rather precise characterization of mso definable string transductions in terms of Hennie transductions, both in the deterministic and in the nondeterministic case. Intuitively an important reason for this equivalence is the inherent boundedness of both types of transductions: mso definable transductions have a bound on the number of copies, whereas Hennie machines have a bound on the number of visits to each of the tape positions.

In case of determinism these two families are equal to the family of transductions realized by two-way generalized sequential machines, Theorem 23. This should be contrasted to nondeterministic transductions, where 2gsm are unable to record choices made during the computation, whereas Hennie machines may use their tape for this purpose.

We summarize.

48 Theorem.

1. MSOS = DHM = 2DGSM.
2. NMSOS = NHM = MREL \circ 2DGSM = 2NGSM_{fin}².

Now that the families NMSOS and 2NGSM have shown to be incomparable, unlike their deterministic counterparts, one may look for natural variants of the families that have the same power. For machines we have discussed such a variant. Indeed, by extending the model with the power of rewriting its input tape (and at the same time demanding the finite visit property) we obtain the Hennie transductions. We leave it as an open problem how to introduce a variant of nondeterminism for mso definable transductions that corresponds to 2ngsm. Additionally, we did not consider transductions realized by one-way transducers. Another remaining problem of interest is the power of first-order logic to define string transductions (where, in Definition 10, we assume all formulas to be first-order, see Example 14). Note that even for $C = \{1\}$ there are first-order definable string transductions that cannot be realized by one-way transducers (such as transforming a string into its reversal). The class of first-order definable string transductions (with respect to nd-gr) such that $C = \{1\}$ and $\phi_*^{1,1}(x, y) = \text{edge}_*(x, y)$ is characterized in [LMSV] to be the class of all transductions that can be realized by functional aperiodic nondeterministic one-way sequential machines (where a sequential machine is a gsm that outputs exactly one symbol at each step). The equivalence of aperiodic finite state automata and first-order logic was established in [MNP71].

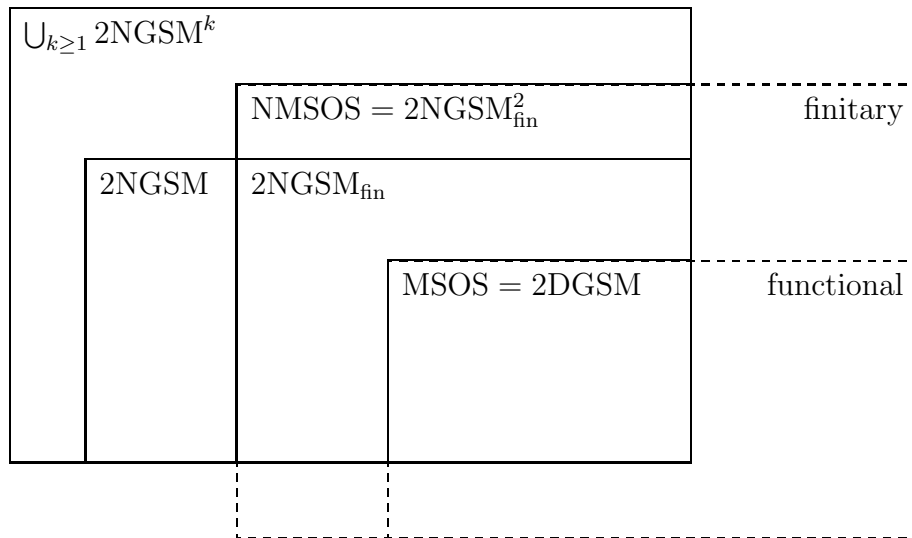


Figure 10: Relationships between our main families of transductions

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